

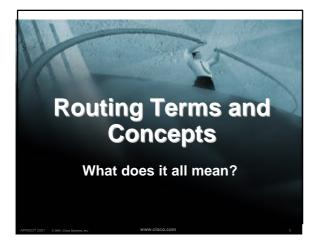
## Introduction • Presenter: Philip Smith, Consulting Engineer Office of the CTO, Cisco Systems e-mail: pfs@cisco.com • Please ask questions

## Agenda

- Routing Terms and Concepts
- Introduction to IGPs
- BGP for ISPs
- Routing Design for ISPs
- Routing Etiquette and the IRR

### Goals

- Promoting a healthy Internet
- Efficient and Effective Routing Configuration
- Internet Routing Registry awareness understanding participation



# Routed backbone HDLC or PPP links between routers Easier routing configuration and debugging Switched backbone Frame Relay/ATM switches in core Surrounded by routers Complex routing & debugging Traffic Engineering

### **PoP Topologies**

- Core routers high speed trunk connections
- Distribution routers and Access routers high port density
- Border routers connections to other AS's
- · Service routers hosting and servers
- Some functions might be handled by a single router

### **Transit, Peering and Default**

- Transit carrying traffic across a network, usually for a fee
- Peering exchanging routing information and traffic
- Default where to send traffic when there is no explicit match is in the routing table

Peering and Transit example

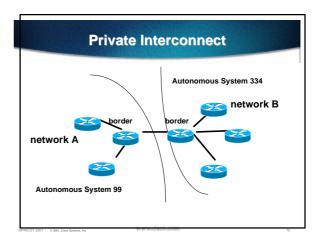
provider A

| IXP-West |
| Provider D

| Provider D

| Provider D

| Provider C |
| A and B can peer, but need transit arrangements with D to get packets



### **Public Interconnect Points**

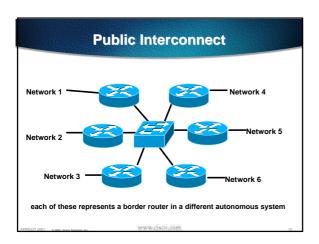
• IXP - Internet eXchange Point

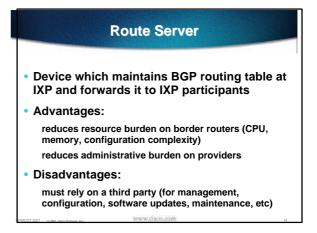
to/from C

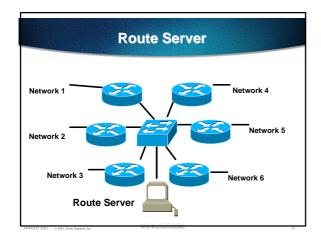
- NAP Network Access Point
- local IXPs
   peering point for a group of local/regional providers
- transit IXPs
   connects local providers to backbone (transit) providers
- hybrid IXPs
   combines the function of local and transit

### **Public Interconnect Point**

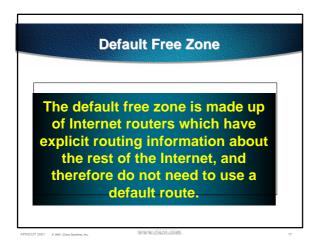
- Centralised (in one facility)
- Distributed (connected via WAN links)
- Shared, switched or routed interconnect Router, FDDI, Ethernet, ATM, Frame relay, SMDS, etc.
- Each provider establishes relationship with other provider at IXP
  - ISP border router peers with all other provider border routers

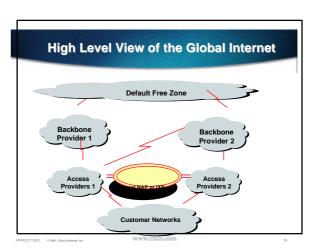


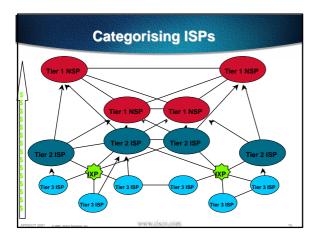












### Inter-provider relationships

- Peering between equivalent sizes of service providers (e.g. Tier 2 to Tier 2)
  - shared cost private interconnection, equal traffic flows
  - "no cost peering"
- Peering across exchange points
   if convenient, of mutual benefit, technically
   foscible
- Fee based peering unequal traffic flows, "market position"



### **IP Addressing**

- Internet is classless
- Concept of Class A, class B or class C is no more
  - engineers talk in terms of prefix length, for example the class B 158.43 is now called 158.43/16.
- All routers must be CIDR capable
  - Classless InterDomain Routing
  - **RFC1812 Router Requirements**

### **IP Addressing**

- Pre-CIDR (<1994)</li>
  - big networks got a class A medium networks got a class B small networks got a class C
- Nowadays
  - allocations/assignments made according to demonstrated need CLASSLESS

**IP Addressing** 

- IPv4 Address space is a resource shared amongst all Internet users
  - Regional Internet Registries delegated allocation responsibility by the IANA
  - APNIC, ARIN, RIPE NCC are the three RIRs
  - RIRs allocate address space to ISPs and Local Internet Registries
  - ISPs/LIRs assign address space to end customers or other ISPs
- 51% of available IPv4 address space used

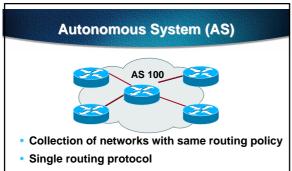
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### **Definitions**

- Non-portable 'provider aggregatable' (PA)
  - Customer uses RIR member's address space while connected to Internet
  - Customer has to renumber to change ISP Aids control of size of Internet routing table
- May fragment provider block when multihoming
- PA space is allocated to the RIR member with the requirement that all assignments are announced as an aggregate

### **Definitions**

- Portable 'provider independent' (PI)
  - Customer gets or has address space independent of ISP
  - Customer keeps addresses when changing ISP
  - Bad for size of Internet routing table
  - PI space is rarely distributed by the RIRs



- Usually under single ownership, trust and administrative control
- AS number obtained from RIR or upstream ISP



### **Routing versus Forwarding**

- Routing = building maps and giving directions
- Forwarding = moving packets between interfaces according to the "directions"



### IP Routing - finding the path

- Path derived from information received from a routing protocol
- · Several alternative paths may exist best next hop stored in forwarding table
- Decisions are updated periodically or as topology changes (event driven)
- Decisions are based on:

topology, policies and metrics (hop count, filtering, delay, bandwidth, etc.)

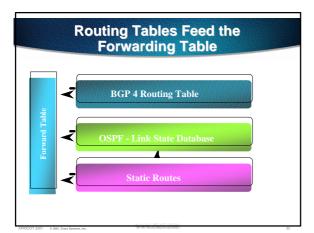
### **IP** route lookup

- Based on destination IP packet
- "longest match" routing more specific prefix preferred over less specific prefix

example: packet with destination of 10.1.1.1/32 is sent to the router announcing 10.1/16 rather than the router announcing 10/8.

### **IP Forwarding**

- Router makes decision on which interface a packet is sent to
- Forwarding table populated by routing process
- Forwarding decisions:
  - destination address
  - class of service (fair queuing, precedence, others) local requirements (packet filtering)
- Can be aided by special hardware



### **Explicit versus Default routing**

- Default:
  - simple, cheap (cycles, memory, bandwidth) low granularity (metric games)
- Explicit (default free zone)
   high overhead, complex, high cost, high granularity
- Hybrid
  - minimise overhead provide useful granularity requires some filtering knowledge

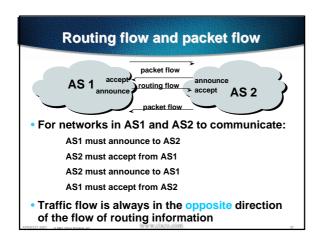
### **Egress Traffic**

- How packets leave your network
- · Egress traffic depends on:
  - route availability (what others send you) route acceptance (what you accept from others)
  - policy and tuning (what you do with routes from others)

Peering and transit agreements

### **Ingress Traffic**

- How packets get to your network and your customers' networks
- Ingress traffic depends on:
   what information you send and to whom
   based on your addressing and AS's
   based on others' policy (what they accept
   from you and what they do with it)



### What Is an IGP?

- Interior Gateway Protocol
- Within an Autonomous System
- Carries information about internal prefixes
- Examples OSPF, ISIS, EIGRP...

terante channe come.

### What Is an EGP?

- Exterior Gateway Protocol
- Used to convey routing information between Autonomous Systems
- De-coupled from the IGP
- Current EGP is BGP4

### Why Do We Need an EGP?

- Scaling to large network
   Hierarchy
   Limit scope of failure
- Policy

Control reachability to prefixes

Merge separate organizations

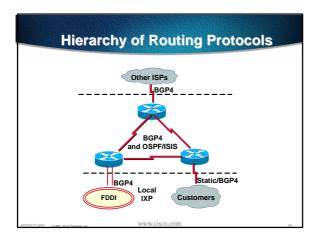
Connect multiple IGPs

### Interior versus Exterior Routing Protocols

- Interior
  - automatic neighbour discovery
  - generally trust your IGP routers
  - routes go to all IGP routers
  - binds routers in one AS together
- Exterior
  - specifically configured peers
  - connecting with outside networks
  - set administrative
  - boundaries
  - binds AS's together
- ADDICOT COOL .....

## Interior versus Exterior Routing Protocols

- Interior
  - Carries ISP infrastructure addresses only
  - ISPs aim to keep the IGP small for efficiency and scalability
- Exterior
  - Carries customer prefixes
  - Carries Internet prefixes
  - EGPs are independent of ISP network topology
- discourant.





### ISIS - Intermediate System to Intermediate System

- Link State Routing Protocol
- OSI development now continued in IETF
- Supports VLSM
- Low bandwidth requirements
- Supports two levels

The backbone (level 2) and areas (level 1)

Route summarisation

### **OSPF - Open Shortest Path First**

- Link State Routing Protocol
- Designed by IETF for TCP/IP RFC2328
- Supports VLSM
- Low bandwidth requirements
- Supports different types of areas
- Route summarisation and authentication

# Why Areas - OSPF Example Backbone Area #0 Area #1 Area #2 Area #3 • Topology of an area is invisible from outside of the area • Results in marked reduction in routing traffic

### Scalable Network Design

• ISIS

Implement level1 - level 2/level 1 hierarchy for large networks only

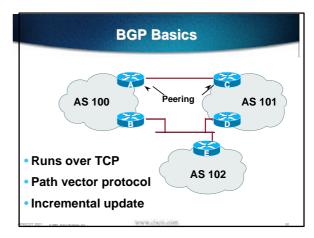
Internet friendly enhanced features

OSPF

Implement area hierarchy
Enforces good network design

- Requires Addressing Plan
- Implement Route Summarisation

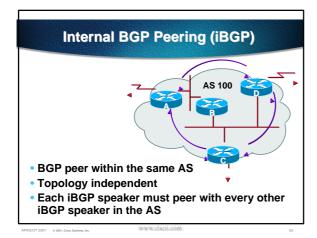


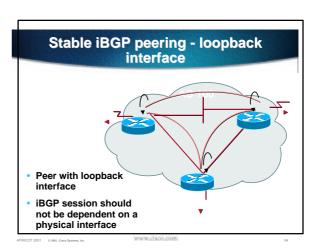


### **BGP General Operation**

- Learns multiple paths via internal and external BGP speakers
- Picks the best path and installs in the IP forwarding table
- Policies applied by influencing the best path selection

# External BGP Peering (eBGP) AS 100 Between BGP speakers in different AS Should be directly connected





### **BGP Attributes**

- Describes characteristics of a prefix
- Some BGP attributes:

AS path, Next hop, Local preference, Multi-Exit Discriminator (MED), Origin, Aggregator and Community.

Some are mandatory, some are transitive

### **BGP Path Selection Algorithm**

- Do not consider path if no route to next hop
- Highest local preference (global within AS)
- Shortest AS path
- Lowest origin code
   IGP < EGP < incomplete</li>

## BGP Path Selection Algorithm (continued)

- Multi-Exit Discriminator
   Considered only if paths are from the same AS
- Prefer eBGP path over iBGP path
- Path with shortest next-hop metric wins
- Lowest router-id

### **BGP in ISP Backbones**

- All routers take part in BGP
- BGP are used to carry some or all of the Internet routing table customer prefixes
- IGPs are used to carry next hop and internal network information recursive route lookup
- Routes are never redistributed from BGP into the IGP or from the IGP into BGP

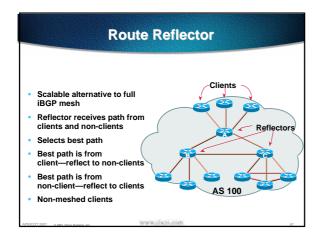


### **Scaling Techniques**

- Administrative scaling (BGP Communities)
- Router resource scaling Route Reflectors (Confederations) Route Flap Damping Dynamic Reconfiguration

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## **Route Reflector**

- Divide the backbone into multiple clusters (hint - build on OSPF/ISIS areas)
- At least one route reflector and few clients per cluster
- · Route reflectors are fully meshed
- Clients in a cluster could be fully meshed
- Single IGP to carry next hop and local routes

### **Route Reflector: Benefits**

- Solves iBGP mesh problem
- Packet forwarding is not affected
- Normal BGP speakers co-exist
- Multiple reflectors for redundancy
- Easy migration
- Multiple levels of route reflectors

## **Route Reflector: Migration** AS 300 AS 100 AS 200 Migrate small parts of the network, one part at a time.



### **Route Flap Damping**

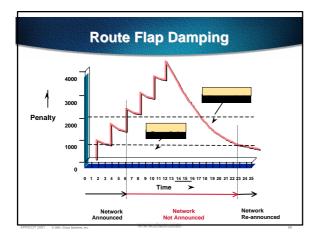
- Route flap
  - Going up and down of path or change in attribute BGP WITHDRAW followed by UPDATE = 1 flap eBGP neighbour going down/up is NOT a flap Ripples through the entire Internet
  - Wastes CPU
- Damping aims to reduce scope of route flap propagation

## Route Flap Damping (Continued)

- Requirements
  - Fast convergence for normal route changes History predicts future behaviour Suppress oscillating routes Advertise stable routes
- Implementation described in RFC2439

### Route Flap Damping Operation

- · Add penalty (1000) for each flap
- Exponentially decay penalty half life determines decay rate
- Penalty above suppress-limit do not advertise route to BGP peers
- Penalty decayed below reuse-limit re-advertise route to BGP peers



## Route Flap Damping Operation

- Only applied to inbound announcements from eBGP peers
- Alternate paths still usable
- In Cisco IOS, controlled by:

Half-life (default 15 minutes)
reuse-limit (default 750)
suppress-limit (default 2000)
maximum suppress time (default 30 minutes)

## Route Flap Damping Configuration

• Examples - ×

bgp dampening 30 750 3000 60

reuse-limit of 750 means maximum possible penalty is 3000 - no prefixes suppressed as penalty cannot exceed suppress-limit

Examples - ✓

bgp dampening 30 2000 3000 60

reuse-limit of 2000 means maximum possible penalty is 8000 - suppress limit is easily reached

### Flap Dampening: Enhancements

- Selective dampening based on AS-path, Community, Prefix
- Variable dampening recommendations for ISPs

http://www.ripe.net/docs/ripe-210.html

Flap statistics in Cisco IOS

show ip bgp neighbor < x.x.x.x> [dampened-routes | flap-statistics]

.....

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### **Soft Reconfiguration**

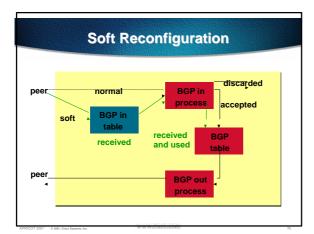
### Problem:

- Hard BGP peer clear required after every policy change because the router does not store prefixes that are denied by a filter
- Hard BGP peer clearing consumes CPU and affects connectivity for all networks

### Solution:

Soft-reconfiguration

Seminar Harris No. 201496



### **Soft Reconfiguration**

- New policy is activated without tearing down and restarting the peering session
- · Per-neighbour basis
- Use more memory to keep prefixes whose attributes have been changed or have not been accepted

### **Configuring Soft reconfiguration**

router bgp 100
neighbor 1.1.1.1 remote-as 101
neighbor 1.1.1.1 route-map infilter in
neighbor 1.1.1.1 soft-reconfiguration inbound
! Outbound does not need to be configured!
Then when we change the policy, we issue an exec command

### **Route Refresh Capability**

- Facilitates non-disruptive policy changes
- No configuration is needed
- No additional memory is used
- Requires peering routers to support "route refresh capability" - RFC2842
- clear ip bgp x.x.x.x in tells peer to resend full BGP announcement

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clear ip bgp 1.1.1.1 soft [in | out]

## Soft Reconfiguration vs Route Refresh

- Use Route Refresh capability if supported find out from "show ip bgp neighbor" does not require additional memory
- Otherwise use Soft Reconfiguration
- Be nice to the Internet!

New York Control of



### **Address Space**

- Approach upstream ISP or consider RIR membership for address space
- Supply addressing plan when requested remember Internet is classless addresses assigned according to need not want
- Assign addresses to backbone and other network layers - remember scalability!
- Some examples follow…

### **Principles of Addressing**

 Separate customer & infrastructure address pools

Manageability

Different personnel manage infrastructure and assignments to customers

Scalability

Easier renumbering - customers are difficult, infrastructure is easy

### **Principles of Addressing**

 Further separate infrastructure In the IGP:

p2p addresses of backbone connections router loopback addresses

Not in the IGP:

RAS server address pools Virtual web and content hosting LANs Mail, DNS servers

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### **Principles of Addressing**

Customer networks

Carry in iBGP

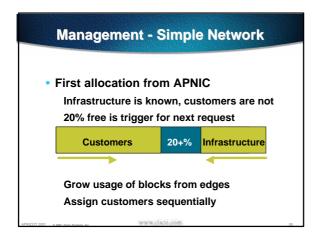
Do not put in IGP – ever!

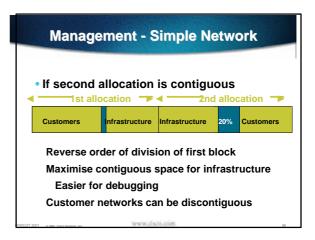
 Do not need to aggregate customer assigned address space

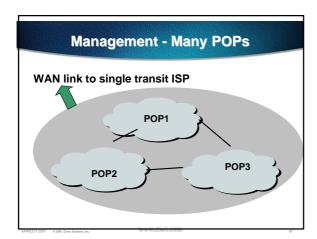
iBGP can carry in excess of 200,000 prefixes, no IGP is designed to do this

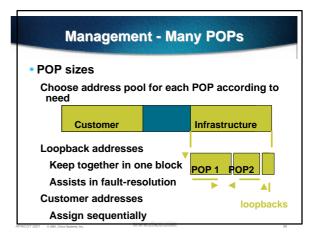
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Management - Many POPs

- /20 minimum allocation is not enough for all your POPs?

Deploy addresses on infrastructure first

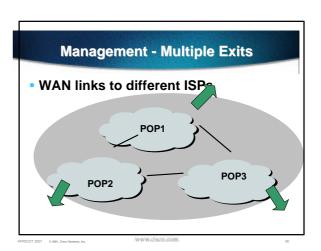
- Common mistake:

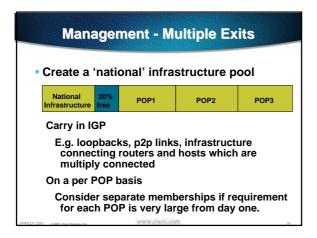
Reserving customer addresses on a per POP basis

- Do not constrain network plans due to lack of address space

Re-apply once address space has been used

There is plenty of it!







### **Network Design**

- Aim for simplicity, scalability and reliability
- Plan the network coverage
- · Estimate growth over the next year
- Design the network

### **Network Coverage**

- Where will you start and how?
- Where will it grow?
   One year is a long time in the Internet
   Future PoP sites
- How big will it grow?
   Inter-site bandwidth availability
- Does it match the business plan?

### **Network Design**

- Start as you mean to continue
- Design scalability from day one hierarchy separate functions
- Choose your IGP carefully scalability, standards knowledge and expertise

### **Designed in Redundancy**

- Design goal should be two of everything
  - Each site should have at least two backbone WAN connections
  - Consider two core routers for each backbone site
- Out of Band management network
- Test lab/network
- Documentation!

### **Deploying IGP**

- Keep IGP small!
  - Smaller IGP, faster convergence in case of link problems
  - Use BGP for customer prefixes, dial pools, and other networks
- Use summarisation between areas of network hierarchy
- Use ip unnumbered where possible

### **External Connections**

- Don't need BGP from day one apply for an AS and deploy BGP only when it is needed i.e. when multihoming
- When deploying BGP
   iBGP carries customer networks only
   IGP carries network link information only
   Do not distribute BGP routes into IGP

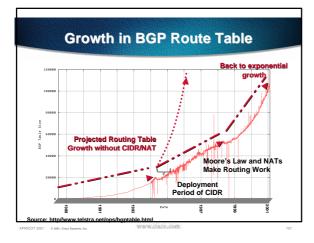
Routing Etiquette

Being a good Internet citizen

### "Problems on the Internet"

- Concern about rate of Internet growth http://www.isc.org/ds/
- Large number of routes
   http://www.employees.org/~tbates/cidr.plot.html
- Routing instability
   http://www.merit.edu/ipma/reports
- Difficulties diagnosing problems
- Quality of Service??

and vice-versa



### **Effects of CIDR on Internet**

- Currently around 100000 routes
   Still too big
- If Internet were unaggregated Would be over 300000 networks (?) May have run out of IPv4 addresses What size of routers required? How stable would the Internet be?

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### **CIDR - Examples**

- Must announce network block assigned by RIR or upstream ISP
- Do not announce subnets of network block, or subnets of other ISPs' network blocks unless exceptional circumstances
- On Cisco routers use redistribute static, or aggregate-address, or network/mask pair

**CIDR - Examples** router bgp 1849 network 194.216.0.0 Redistribute redistribute static static ! Must have a matching IGP route ip route 194.216.0.0 255.255.0.0 null0 router bgp 1849 Aggregate network 194.216.0.0 aggregate-address 194.216.0.0 255.255.0.0 address ! More specific route must exist in BGP table router bgp 1849 network 194.216.0.0 mask 255.255.0.0 Network/mask pair ! Must have a matching IGP route ip route 194.216.0.0 255.255.0.0 null0

### **CIDR - Positive Efforts**

- Most ISPs now filter all prefixes longer than /24
- Some ISPs pay attention to Tony Bates' CIDR report
- Some ISPs filter according to policy registered in the Internet Routing Registry
- No aggregation or bad aggregation could result in no connectivity

### Aggregation

- Announce aggregate to rest of Internet
- Put it into Routing Registry (route object)
- Keep more specifics internal to network
   Use iBGP for carrying customer networks
   Use IGP for carrying backbone addresses
   Aggregate internally when possible

### **Aggregation - Good Example**

- Customer link goes down their /26 network becomes unreachable
- /19 aggregate is still being announced no BGP hold down problems no BGP propagation delays no dampening by other ISPs

**Aggregation - Good Example** 

- Customer link returns
- Their /26 network is visible again
- The whole Internet becomes visible immediately
- Quality of Service perception

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### **Aggregation - Bad Example**

- · Customer link goes down
  - Their /23 network becomes unreachable
- Their ISP doesn't aggregate their /19 network block
  - /23 network withdrawal announced to peers starts rippling through the Internet
  - added load on all Internet backbone routers as network is removed from routing table

### **Aggregation - Bad Example**

- Customer link returns
  - Their /23 network is now visible to their ISP
  - Their /23 network is re-advertised to peers
  - Starts rippling through Internet
  - Load on Internet backbone routers as network is reinserted into routing table
  - Some ISP's dampen flaps
  - Internet may take 10-20 min or longer to be visible
  - Quality of Service???

### **Aggregation - Summary**

- Good example is what everyone should do!
  - Adds to Internet stability
  - Reduces size of routing table
  - Reduces routing churn
  - Improves Internet QoS for everyone
- Bad example is what many still do!
  - Laziness? Lack of knowledge?

### "The New Swamp" - Feb 2001

- Areas of poor aggregation
- 192/3 space contributes 78000 networks rest of Internet contributes 22000 networks

	<del>letworks</del>	Elock	Networks	Block	Networks	Block	Networks
192/8	6602	200/8	2902	208/8	4987	217/8	400
193/8	2908	201/8		209/8	5392	24/8	1466
194/8	3122	202/8	4174	210/8	1445	61/8	230
195/8	1839	203/8	7280	211/8	882	62/8	575
196/8	604	204/8	5023	212/8	2193	63/8	2833
197/8	0	205/8	3395	213/8	1049	64/8	3423
198/8	4853	206/8	4523	214/7	23	65/8	283
199/8	4462	207/8	4583	216/8	5391	66/8	470

### "The New Swamp" - July 2000

 192/3 space contributes 69000 networks - rest of Internet contributes 16000 networks

Block	Networks	Block	Networks	Block	Networks	Block	Networks
192/8	6352	200/8	2436	208/8	4804	12/8	1047
193/8	2746	201/8		209/8	4755	24/8	1122
194/8	2963	202/8	3712	210/8	1375	61/8	80
195/8	1689	203/8	5494	211/8	532	62/8	428
196/8	525	204/8	4694	212/8	1859	63/8	2198
197/8	0	205/8	3210	213/8	635	64/8	1439
198/8	4481	206/8	4206	214/7	14		
199/8	4084	207/8	3943	216/8	4177		

13310 4004 20170 3343 2107

### **Original Swamp Cause**

- Early growth of Internet
- Classful network allocation
- Small number of connected networks
- · Lack of foresight by all

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### **New Swamp Persists**

- Lazy or technically naïve ISPs announcing 32 /24s rather than /19 aggregate block
  - announcing customer prefixes as they connect rather than aggregate block only
- Poorly thought out multihoming
- Technical solutions keep ahead of problem so far:

faster routers, more memory, CIDR

### **Solutions**

- Don't route other ISP's address space unless in failure mode during multihoming
- Aggregate!
- Don't announce subprefixes of your assigned block
- Be prudent when announcing small prefixes out of former A and B space

### **Solutions**

- Encourage other ISPs to be good citizens don't route their bad citizenship
- Multihoming fragments address space think carefully about set up and requirements load balancing versus resilience

http://infopage.cw.net/Routing

### **Efforts**

- Tony Bates' CIDR report sent to nanog, apops and eof mail lists
- Routing Report sent to APOPS, ARIN rtma and RIPE routing-wg
- Regional Internet Registries
- Many ISPs still care
- Peer pressure
- YOU!

### **Renumbering - motivation**

- Same as motivation for aggregation holes are bad, using swamp space
- First time Internet connection
   legal address space, practical addressing scheme
- New Provider
   renumber into new provider's block
   reduces fragmentation and improves routability

### Renumbering - how to?

 PIER - Procedures for Internet and Enterprise Renumbering

http://www.isi.edu/div7/pier/papers.html

- Be aware of effect on essential services
   e.g. DNS ttl requires lowering, router filters
- Use DHCP, secondary addressing
- Not difficult but needs planning

### **Route Flap Damping**

- Route Flap technical description earlier
- Many ISPs now suppress route flaps at network borders
- Cisco BGP Case Study at http://www.cisco.com/warp/public/459/16.html
- · Recommended parameters are at

http://www.ripe.net/docs/ripe-210.html

### **Route Flap Damping - Caution**

- Be aware of potential problems
- Unreachability could be due to dampening, not disconnection
- Border routers need more memory and CPU
- Train your staff!

Name for 15th No. 2004

### **Filtering Policies**

- Filter announcements by peers
   AS list, prefixes
- Only accept what is listed in routing registry
  - avoids configuration errors and routing problems authorisation?
- Only announce what you list in routing registry
- · Keep routing registry and filters up to date

### "Documenting Special Use Addresses" - DSUA

 Private and Special Use addresses must be blocked on all BGP peerings, in and out:

http://www.ietf.org/internet-drafts/draft-manning-dsua-06.txt

```
ip prefix-list private-sua deny 0.0.0.0/8 le 32 ip prefix-list private-sua deny 10.0.0.0/8 le 32 ip prefix-list private-sua deny 127.0.0.0/8 le 32 ip prefix-list private-sua deny 127.0.0.0/8 le 32 ip prefix-list private-sua deny 192.16.0.0/12 le 32 ip prefix-list private-sua deny 172.16.0.0/12 le 32 ip prefix-list private-sua deny 192.0.2.0/24 le 32 ip prefix-list private-sua deny 192.168.0.0/16 le 32 ip prefix-list private-sua deny 224.0.0.0/3 le 32 ip prefix-list private-sua deny 0.0.0.0/0 ge 25 ip prefix-list private-sua permit 0.0.0.0/0 le 32
```

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### **Definition**

 "A public authoritative distributed repository of routing information"

Public databases

Distributed repository of information

Have authoritative data

Vendor independent

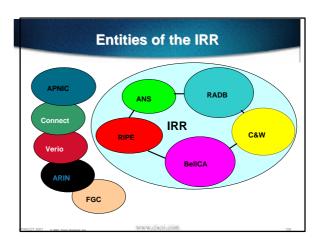
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126

### Composition

- Routing Policy Details
- Routes and their aggregates
- Topology Linking AS's
- Network components such as routers
- Is separate from other information such as domains and networks



### **Relationship Table** Registry APNIC Yes Soon Yes Nο RIPE Yes Yes No RADB Yes Nο Yes Nο C&W Yes ANS Yes No No BellCA Yes ARIN Yes Yes No InterNIC' No

## Relationships 37 RRs around the world C&W, ANS and BellCA - provider run RRs Other RRs run by Verio, FGC, Connect, etc... RIPE RR - European providers ARIN RR - launched 8 February 1999 RADB - Default RR for rest of world APNIC - plans to be full member of IRR very soon.

### Benefits of an IRR

- Operational Support
- Information
- Configuration
- Problem diagnosis
- Improved Service Quality
- Tools for consistency checking

### Information

- Routing policy repository
- "Map of global routing topology"
- Routing policy between neighbouring AS's
- Device independent description of routing policy

### Configuration

- Supports network filtering
- Configures routers and policies
- Revision control
- Sanity checking
- Simulation

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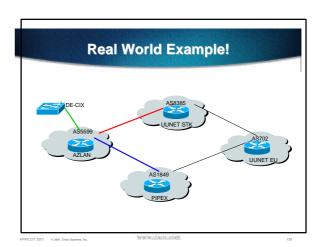
## Key Objects and Syntax of RIPE-181

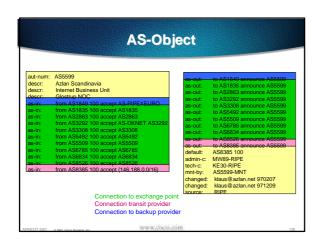
- Representation
- AS Object
- AS Macro
- Route Object
- Authorisation Maintainer Object

Representation

- ASCII printable
- Attributes by tag:value lines
- Objects separated by empty lines
- RIPE-181 and RPSL

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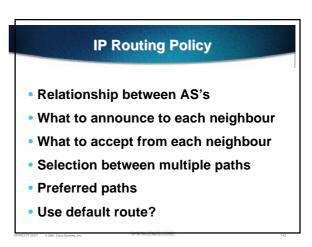


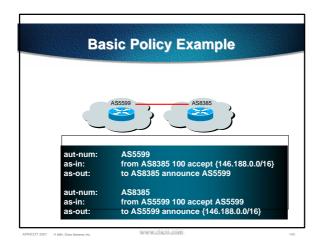


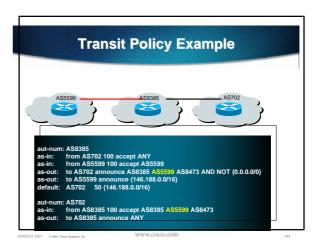
## • Can represent policy using Boolean expressions (AND, OR, NOT) Keyword ANY - means "everything" Communities and AS Macros Route lists - {prefixes} Cost to indicate preference Attribute DEFAULT - accept 0.0.0.0

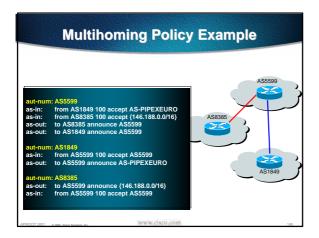
Syntax for AS Object

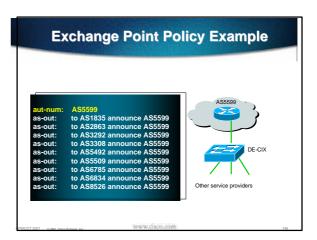
## Mandatory Fields aut-num, descr, admin-c, tech-c, mnt-by, changed, source, as-in, as-out Optional Fields as-name, interas-in, interas-out, as-exclude, default, guardian, remarks, notify









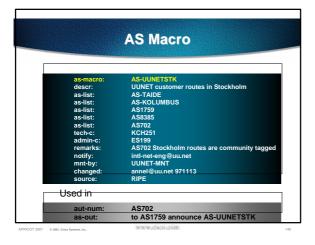


### **AS Macro**

- Collection of AS's or other AS macros
- Describes membership of a set
- · Contains no policy info
- Scales better
- Can differentiate between customer and peer routes

### **Fields in AS Macro**

- Mandatory Fields
  - as-macro, descr, as-list, tech-c, admin-c, mnt-by, changed, source
- Optional Fields
  - guardian, remarks, notify



### **Route Object**

- Represents a route in the Internet
- Contains all membership information
- Only one origin possible
- Classless (should be aggregated)
- Can support holes and withdrawn

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## Fields in Route Object Mandatory Fields route, descr, origin, mnt-by, changed, source Optional Fields hole, withdrawn, comm-list, remarks, notify Example: route: 195.129.0.0/19 descr: UNNET-NET origin: A\$702 remarks: UUNET-filter inbound on prefixes longer than /24 notify: intl-net-eng@uu.net mnt-by: UUNET-MMT changed: annel@uu.net 970501 source: RIPE



## How to register and update information in the IRR

- Frequently used objects
- Update procedures

Modifying Objects

Deleting Objects

Submitting Objects

Authorisation/Notification

**Errors and Warnings** 

**NIC** handles

### **Frequently Used Objects**

- Person contact person
- Maintainer authorisation of objects
- Inetnum address assignment
- Aut-num autonomous systems
- AS-macro set of AS's
- Route announced routes

### **Unique Keys**

- Uniquely identifies an object
- Updating object overwrites old entry need unique key
- Used in querying whois
- Web based full text searches available now, e.g.

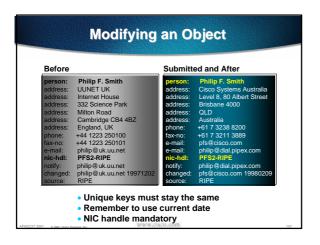
http://whois.apnic.net/apnic-bin/whois.pl

### **Unique Keys**

- Person name plus NIC handle
- Maintainer maintainer name
- Inetnum network number
- Aut-num AS number
- AS-macro AS macro name
- Route route value plus origin

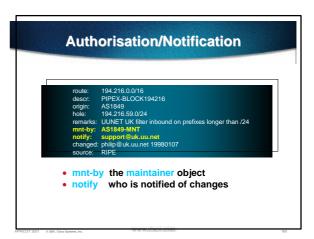
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www.dsco.com.

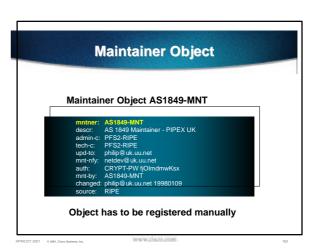


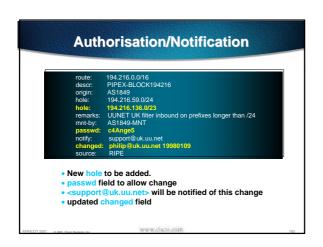


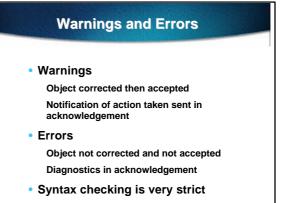
## • Email Interface - eg APNIC auto-dbm@apnic.net Robot mail box Send all database updates to this mailbox Can use LONGACK and HELP in the subject line apnic-dbm@apnic.net human mailbox questions on the database process

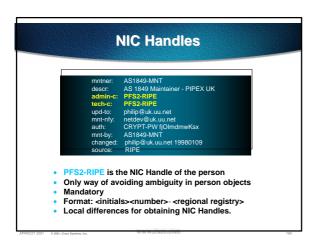


## Maintainer Object Who is authorised Authorisation Method email-from and crypt-pw Mandatory Fields mntner, descr, admin-c, tech-c, upd-to, auth, mnt-by Optional Fields mnt-nfy, changed, notify, source











### What is RPSL

- RPSL is the development of RIPE-181
   RFC2622 Routing Policy Specification Language
- Allows more complex policy specification Looks very similar to RIPE-181 (but not backward compatible)
- All participants in the IRR have agreed to migrate to RPSL

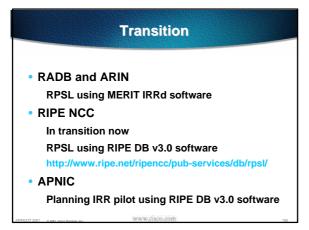
Many already have

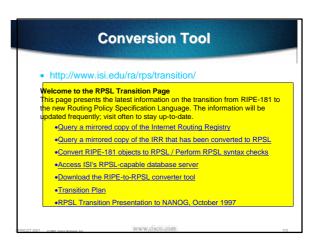
· Training materials at

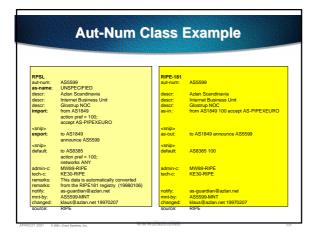
http://www.isi.edu/ra/rps/training

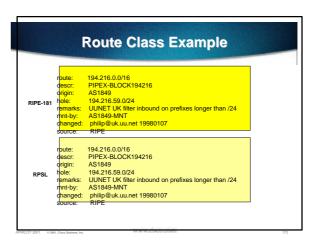
/ww.isi.euu/ra/rps/trailillig

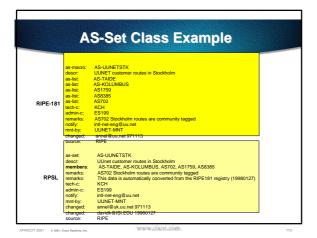
## RPSL Database Software RPSL database software available: IRRd (Merit) – http://www.irrd.net/ In fully deployment RIPE DB v3.0 – http://www.ripe.net/ In beta test RIPE DB 3.0.0b2 with ISI RPSL extensions http://www.kessens.com/~david/software/ BIRD v1.1beta – ftp://ftp.isi.edu/ra/BIRD Status unknown

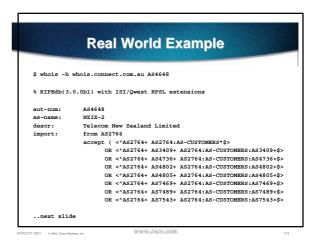


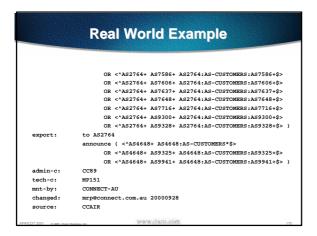














### What tools and resources?

- RAToolset
  - www.isi.edu/ra/RAToolSet
- RIPE whois
  - ftp.ripe.net/ripe/tools
- Traceroute sites
  - www.traceroute.org
- Looking Glasses
  - http://www.traceroute.org/#Looking Glass

### **RAToolSet**

- Two versions
  - 3.5.8 supports RIPE-181
  - 4.6.3 supports RPSL
- Runs on most Unix platforms
- · Requires recent g++, tcl and tk
- Excellent for housekeeping, debugging and configuration

### **RAToolSet Tools**

- RTconfig
  - Generate router configurations for Cisco, Bay, GateD and Juniper
- AOE aut-num object editor update aut-num, as-macro objects
- ROE route-object editor update route-object
- CIDRadvisor
  - advice on CIDRisation

### **ROE Uses**

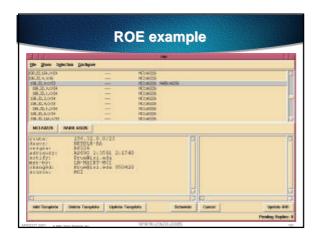
- Route object editor used to:
  - check for consistency of route objects in IRRs
  - synchronise route object entries in different IRRs
  - detect missing or unwanted route objects

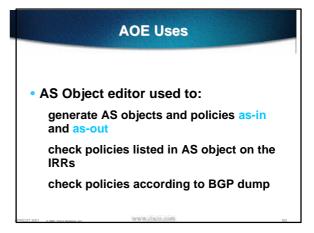
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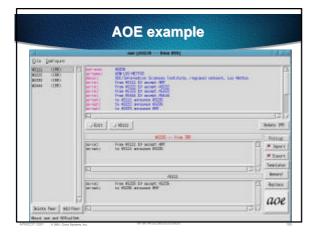
www.dsco.com

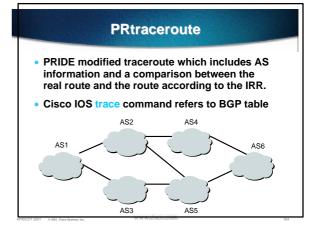
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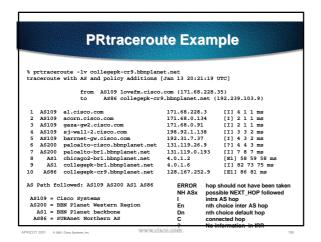
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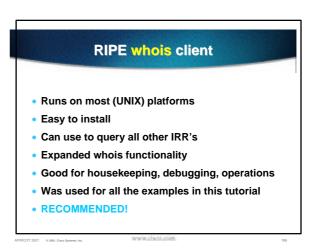




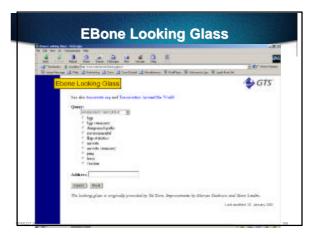




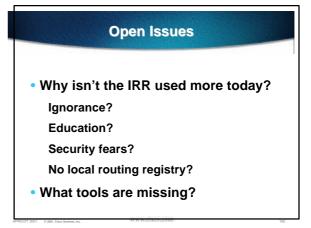




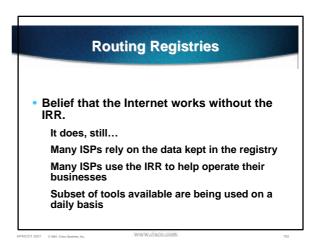








## Should software be available as a commercial package? Better bundled/supported/debugged? Better integration/training? Most tools are freely available public efforts for the good of the "community"

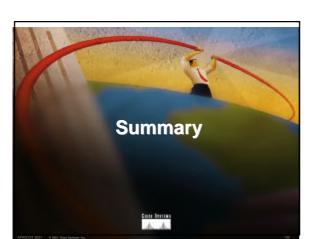


### **Awareness & Training**

- Is there enough awareness about Internet routability?
- Is there enough training on the promotion of routability
- Headcount requirement depends on organisation too easy and cheaper to be irresponsible
- Overall organisational awareness of the issues ® overall efficiency, quality of service and support

### **Ways forward**

- Routing Registry enhancements
   RPSL matches most of BGP's policy capabilities today
- Feedback on tool enhancements
- Feedback to vendors on equipment configuration enhancements
- More training, more education, more feedback!



### **Summary**

- ISP networks and terminology
- The application of IGPs and BGP in an Internet network
- Shown tools which help diagnose and solve routing problems more easily
- Application of routing registries

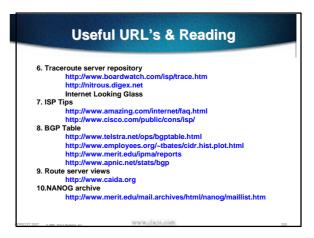
### **Summary**

- Made you more aware of the issues facing the Internet today
- Showed you how to make a positive contribution to the functioning of the Internet
- Promoted Routability!

### The End!

- Any Questions?
- Please fill in evaluation form
- This presentation will be available at http://www.cisco.com/public/cons/isp/documents
- My contact info:
   Philip Smith <pfs@cisco.com>

## 1. CIDR ftp://ftp.isi.edu/in-notes/rfc{1517,1518,1519).txt http://www.ibm.net.ii/-hank/cidr.html ftp://ftp.uninett.no/pub/misc/eidnes-cidr.ps.Z Network addressing when using CIDR 2. AS numbers ftp://ftp.isi.edu/in-notes/rfc1930.txt Guidelines for creation, selection, and registration of an AS 3. Address Allocation and Private Internets ftp://ftp.isi.edu/in-notes/rfc1918.txt 4. BGP Dampening http://www.cisco.com/warp/public/459/16.html ftp://ftp.ripe.net/ripe/docs/ripe-210.txt European recommendations for route flap dampening ftp://engr.ans.net/pub/slides/nanog/feb-1995/route-dampen.ps 5. Routing Discussion http://www.ripe.net/wg/routing/index.html



## 1. RFC1786 "Representation of IP Routing Policies in a Routing Registry" ftp://ftp.isi.edu/in-notes/rfc1786.txt 2. RATools and RSPL ftp://ftp.apnic.net/let/frc/rfc2280.txt Tools http://www.isi.edu/ra/ Mailing List <ratoolset@isi.edu> 3. PRIDE Slides ftp://ftp.ripe.net/pride/docs/course-slides Guide ftp://ftp.ripe.net/pride/docs/guide-2.0txt.{ps}.tar.gz Tools ftp://ftp.ripe.net/pride/tools/\* 4. IRR authorisation/notification ftp://ftp.ripe.net/pride/tools/\* 5. RADB pointers http://www.ra.net http://www.ra.net http://www.ra.net/faq.htm 6. ISP run RR User documents http://infopage.cw.net/Routing