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BGP for Internet Service Providers

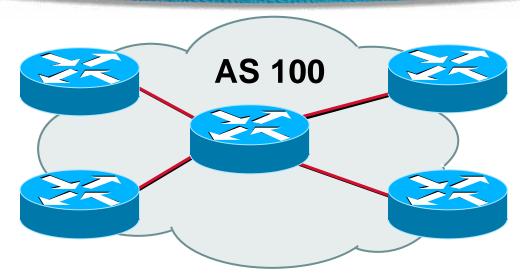
- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

BGP Basics What is this BGP thing? CISCO SYSTEMS NANOG 22 © 2000, Cisco Systems, Inc.

Border Gateway Protocol

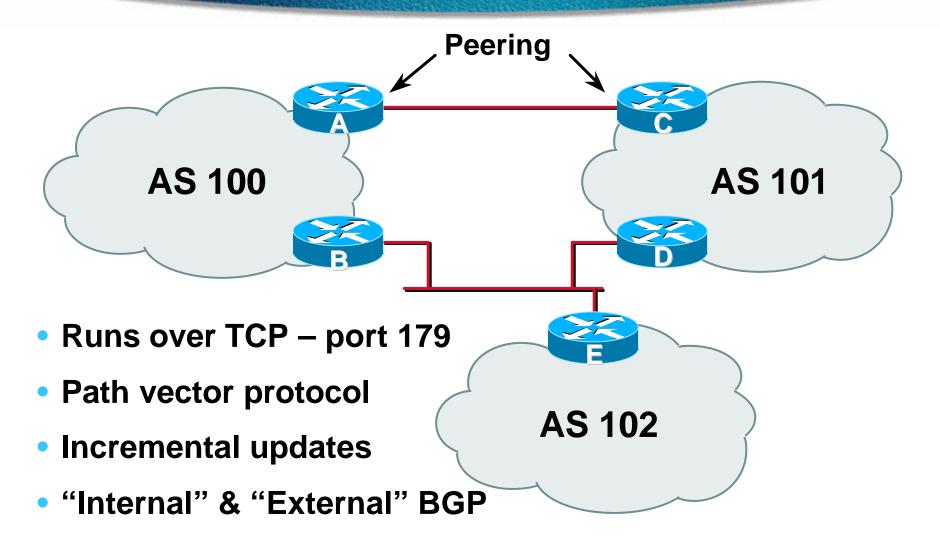
- Routing Protocol used to exchange routing information between networks exterior gateway protocol
- RFC1771
 work in progress to update
 draft-ietf-idr-bgp4-12.txt

Autonomous System (AS)

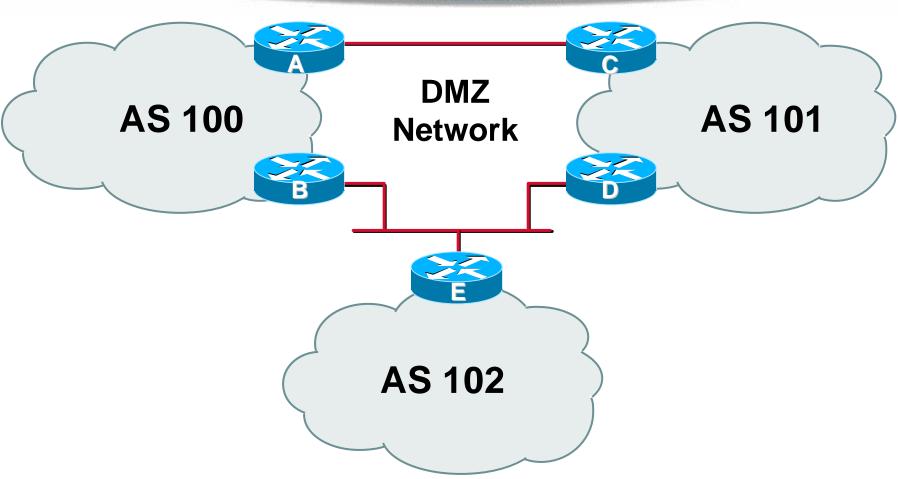


- Collection of networks with same routing policy
- Single routing protocol
- Usually under single ownership, trust and administrative control

BGP Basics



Demarcation Zone (DMZ)

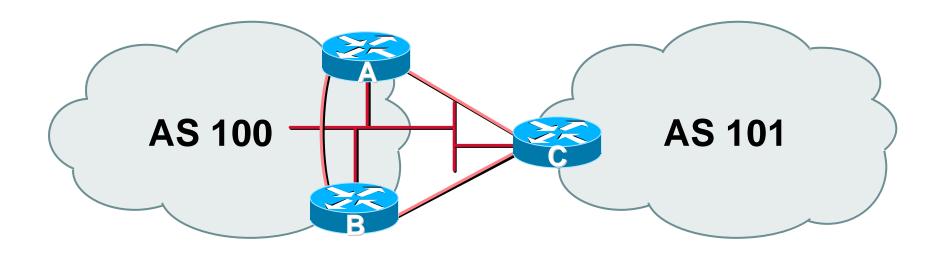


Shared network between ASes

BGP General Operation

- Learns multiple paths via internal and external BGP speakers
- Picks the best path and installs in the forwarding table
- Best path is sent to external BGP neighbours
- Policies applied by influencing the best path selection

External BGP Peering (eBGP)



- Between BGP speakers in different AS
- Should be directly connected
- Never run an IGP between eBGP peers

Configuring External BGP

Router A in AS100

```
interface ethernet 5/0
ip address 222.222.10.2 255.255.255.240
router bgp 100
network 220.220.8.0 mask 255.255.252.0
neighbor 222.222.10.1 remote-as 101
neighbor 222.222.10.1 prefix-list RouterC in neighbor 222.222.10.1 prefix-list RouterC out
```

Router C in AS101

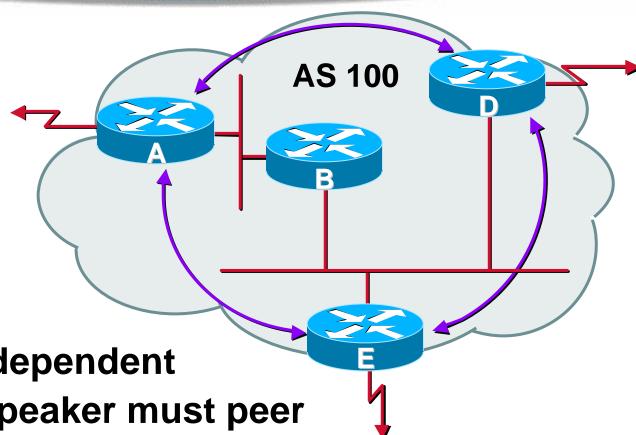
```
interface ethernet 1/0/0
ip address 222.222.10.1 255.255.255.240
router bgp 101
network 220.220.16.0 mask 255.255.240.0
neighbor 222.222.10.2 remote-as 100
neighbor 222.222.10.2 prefix-list RouterA in neighbor 222.222.10.2 prefix-list RouterA out
```

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Internal BGP (iBGP)

- BGP peer within the same AS
- Not required to be directly connected
- iBGP speakers need to be fully meshed they originate connected networks they do not pass on prefixes learned from other iBGP speakers

Internal BGP Peering (iBGP)

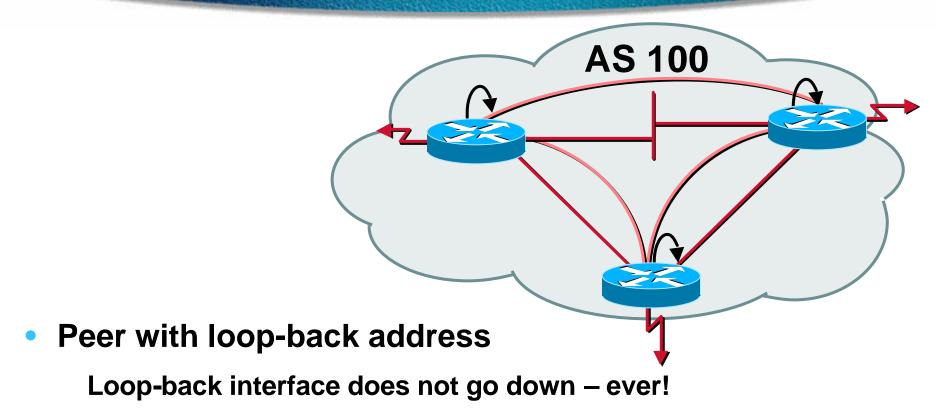


Topology independent

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 Each iBGP speaker must peer with every other iBGP speaker in the AS

Peering to Loop-Back Address



- iBGP session is not dependent on state of a single interface
- iBGP session is not dependent on physical topology

Configuring Internal BGP

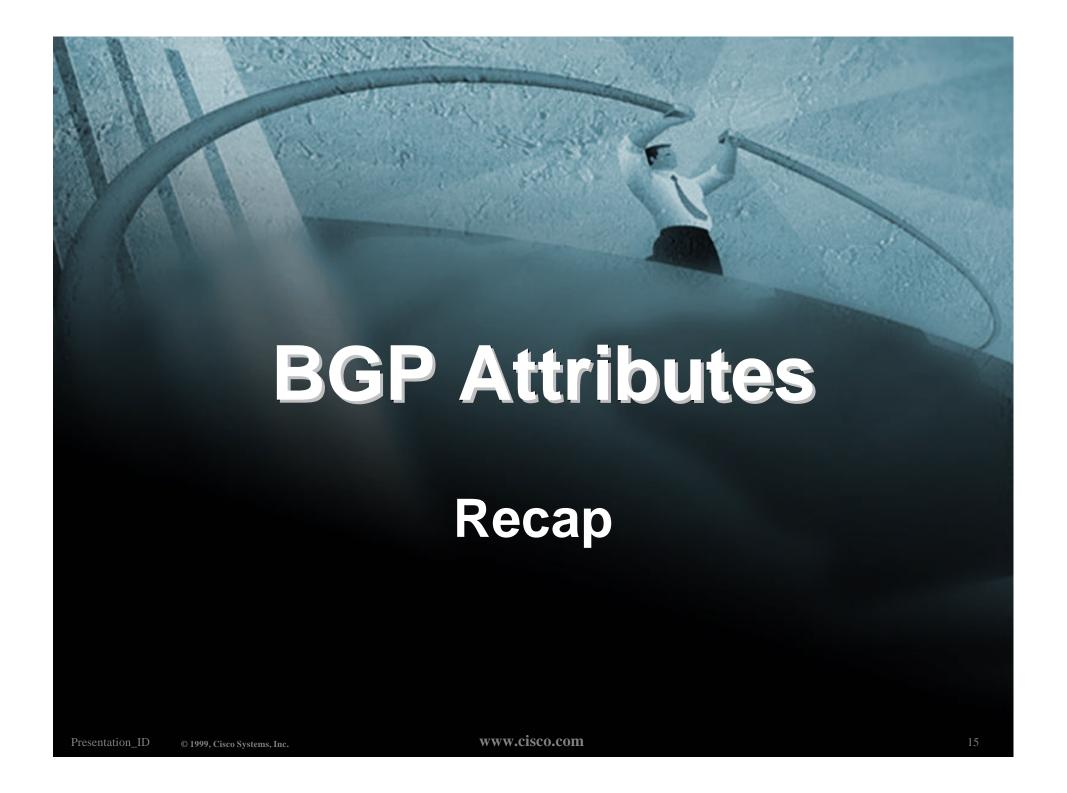
Router A

```
interface loopback 0
ip address 215.10.7.1 255.255.255.255
router bgp 100
  network 220.220.1.0
  neighbor 215.10.7.2 remote-as 100
  neighbor 215.10.7.2 update-source loopback0
  neighbor 215.10.7.3 remote-as 100
  neighbor 215.10.7.3 update-source loopback0
```

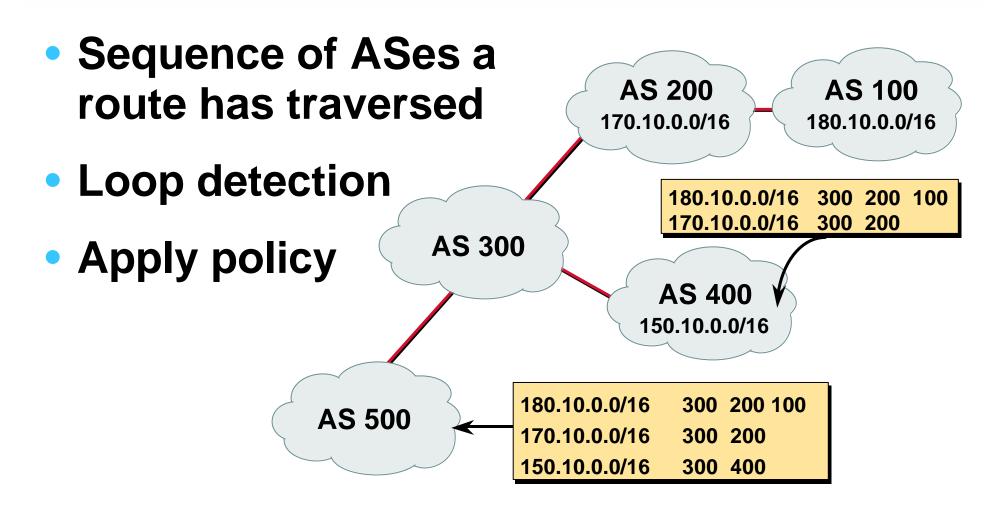
Router B

```
interface loopback 0
ip address 215.10.7.2 255.255.255.255
router bgp 100
  network 220.220.5.0
  neighbor 215.10.7.1 remote-as 100
  neighbor 215.10.7.1 update-source loopback0
  neighbor 215.10.7.3 remote-as 100
  neighbor 215.10.7.3 update-source loopback0
```

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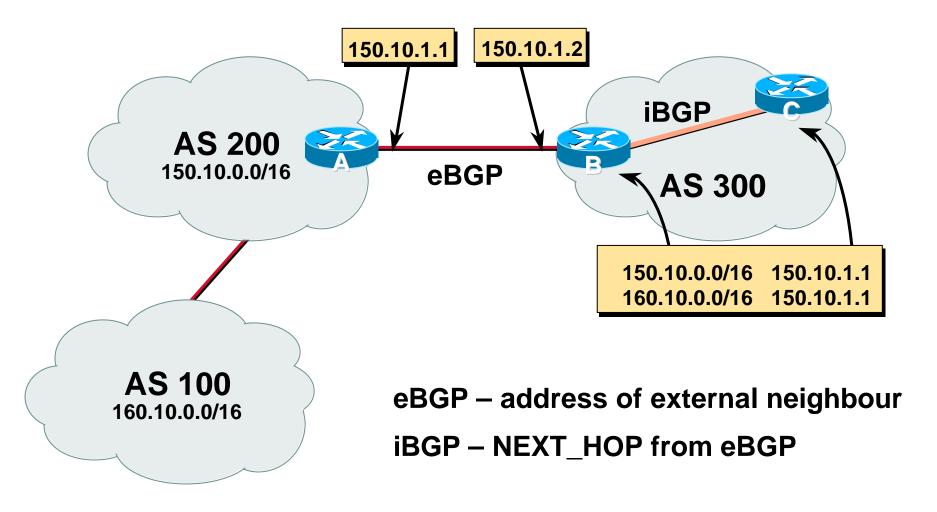


AS-Path



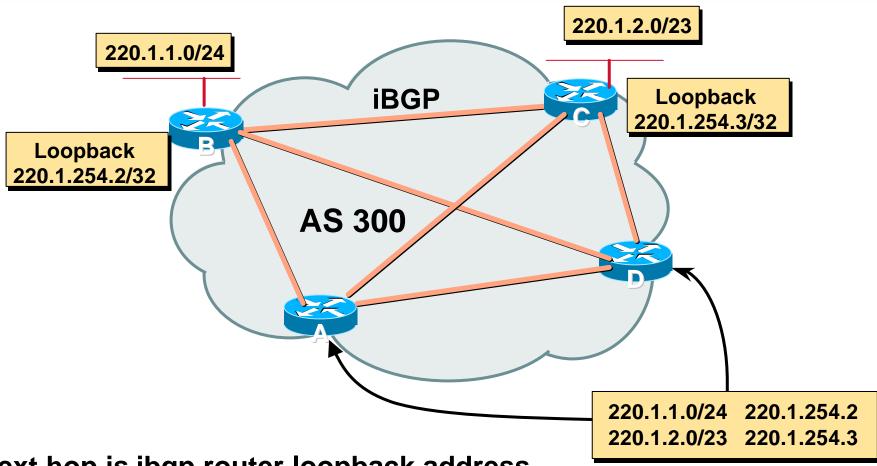
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Next Hop



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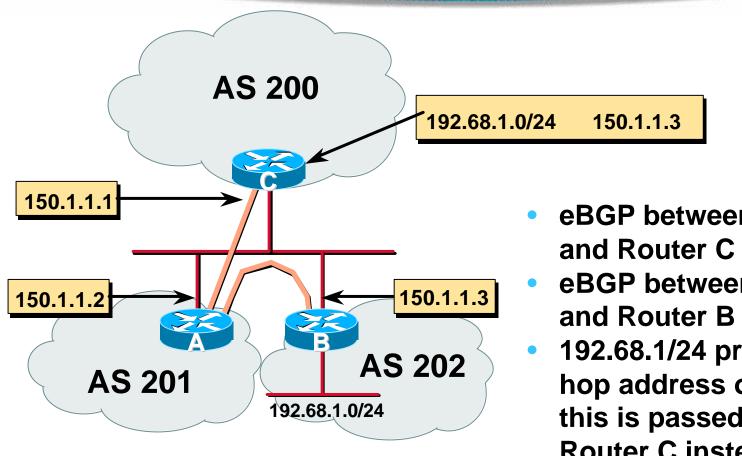
iBGP Next Hop



Next hop is ibgp router loopback address

Recursive route look-up

Third Party Next Hop



- eBGP between Router A
- eBGP between Router A
- 192.68.1/24 prefix has next hop address of 150.1.1.3 this is passed on to Router C instead of 150.1.1.2

Next Hop (summary)

- IGP should carry route to next hops
- Recursive route look-up
- Unlinks BGP from actual physical topology
- Allows IGP to make intelligent forwarding decision

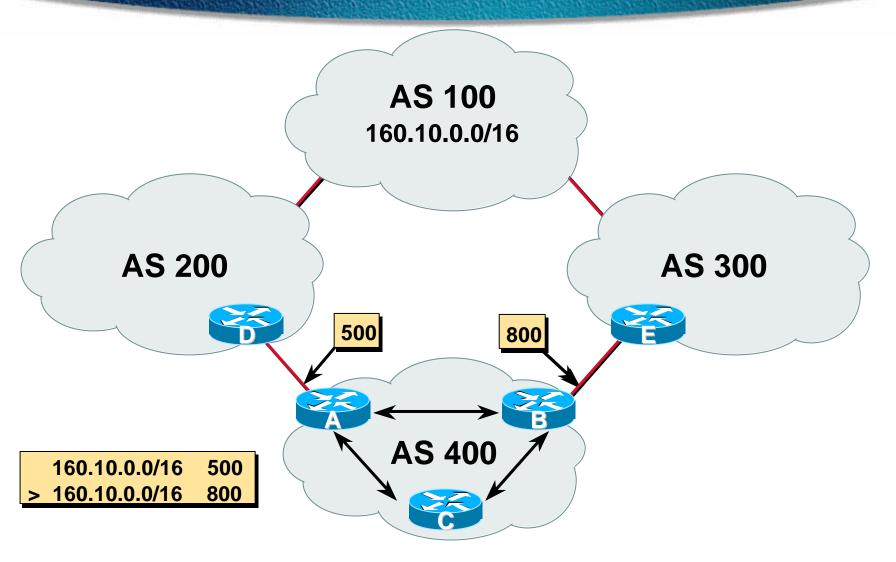
Origin

- Conveys the origin of the prefix
- "Historical" attribute
- Influences best path selection
- Three values: IGP, EGP, incomplete
 - IGP generated by BGP network statement
 - EGP generated by EGP
 - incomplete redistributed from another routing protocol

Aggregator

- Conveys the IP address of the router/BGP speaker generating the aggregate route
- Useful for debugging purposes
- Does not influence best path selection

Local Preference



Local Preference

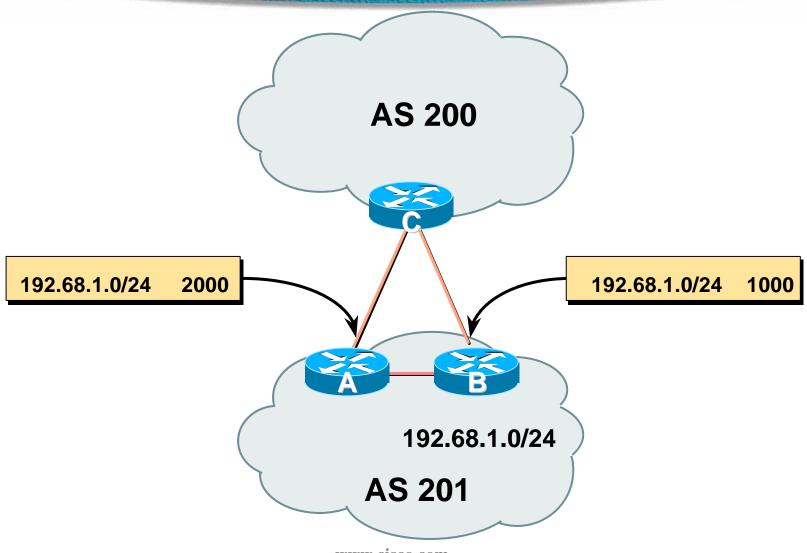
- Local to an AS non-transitive
 Default local preference is 100
- Used to influence BGP path selection determines best path for outbound traffic
- Path with highest local preference wins

Local Preference

Configuration of Router B:

```
router bgp 400
neighbor 220.5.1.1 remote-as 300
neighbor 220.5.1.1 route-map local-pref in
!
route-map local-pref permit 10
match ip address prefix-list MATCH
set local-preference 800
!
ip prefix-list MATCH permit 160.10.0.0/16
```

Multi-Exit Discriminator (MED)



Multi-Exit Discriminator

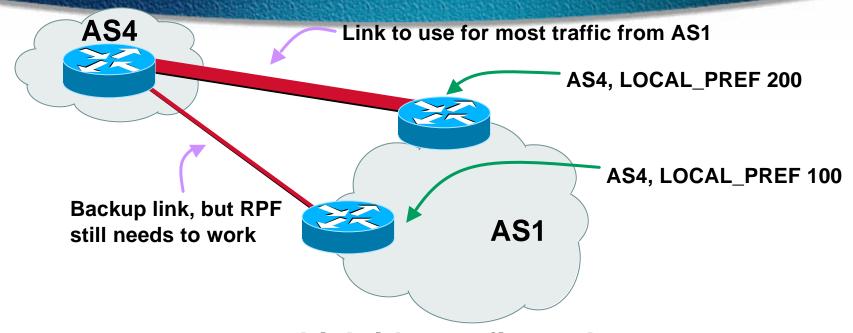
- Inter-AS non-transitive
- Used to convey the relative preference of entry points
 - determines best path for inbound traffic
- Comparable if paths are from same AS
- IGP metric can be conveyed as MED set metric-type internal in route-map

Multi-Exit Discriminator

Configuration of Router B:

```
router bgp 400
neighbor 220.5.1.1 remote-as 200
neighbor 220.5.1.1 route-map set-med out
!
route-map set-med permit 10
match ip address prefix-list MATCH
set metric 1000
!
ip prefix-list MATCH permit 192.68.1.0/24
```

Weight – used to deploy RPF



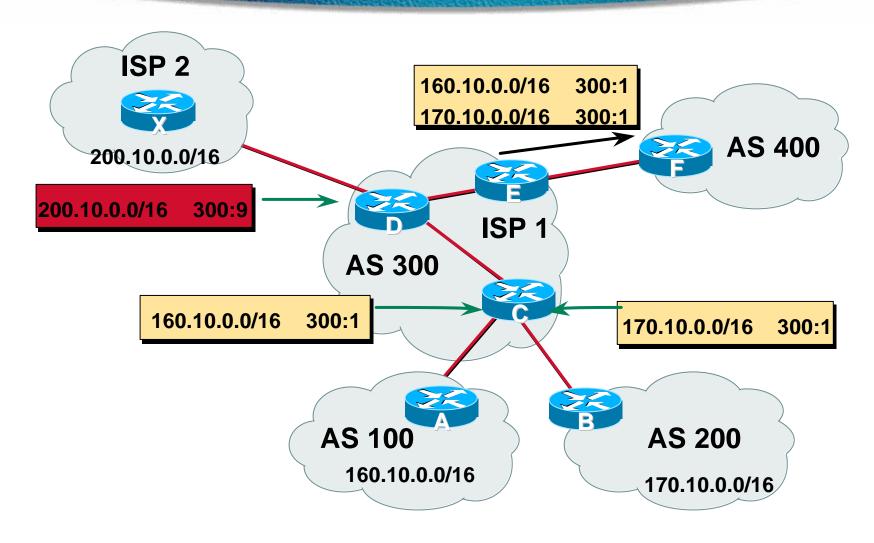
- Local to router on which it's configured
 Not really an attribute
- route-map: set weight
- Highest weight wins over all valid paths
- Weight customer eBGP on edge routers to allow RPF to work correctly

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Community

- BGP attribute
- Described in RFC1997
- 32 bit integer
 Represented as two 16bit integers
- Used to group destinations
 Each destination could be member of multiple communities
- Community attribute carried across AS's
- Very useful in applying policies

Community



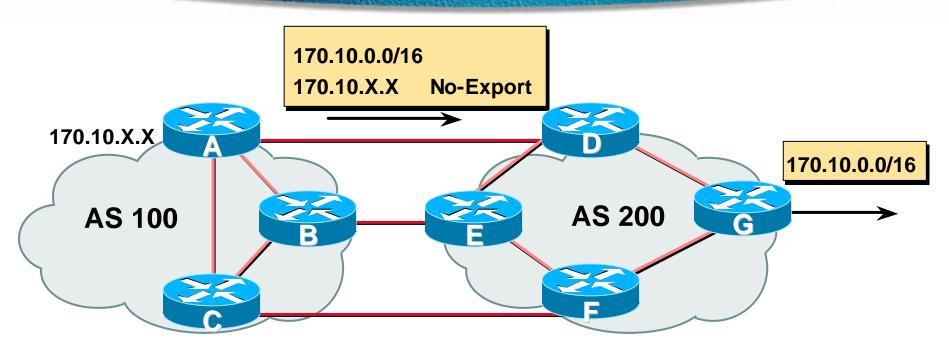
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Well-Known Communities

- no-export
 do not advertise to eBGP peers
- no-advertise
 do not advertise to any peer
- local-AS

do not advertise outside local AS (only used with confederations)

No-Export Community



- AS100 announces aggregate and subprefixes aim is to improve loadsharing by leaking subprefixes
- Subprefixes marked with no-export community
- Router G in AS200 strips out all prefixes with no-export community set



BGP Path Selection Algorithm

- Do not consider path if no route to next hop
- Do not consider iBGP path if not synchronised (Cisco IOS)
- Highest weight (local to router)
- Highest local preference (global within AS)
- Prefer locally originated route
- Shortest AS path

BGP Path Selection Algorithm (continued)

- Lowest origin code
 IGP < EGP < incomplete
- Lowest Multi-Exit Discriminator (MED)

If bgp deterministic-med, order the paths before comparing

If bgp always-compare-med, then compare for all paths

otherwise MED only considered if paths are from the same AS (default)

BGP Path Selection Algorithm (continued)

- Prefer eBGP path over iBGP path
- Path with lowest IGP metric to next-hop
- Lowest router-id (originator-id for reflected routes)
- Shortest Cluster-List

Client must be aware of Route Reflector attributes!

Lowest neighbour IP address



Applying Policy with BGP

Applying Policy

Decisions based on AS path, community or the prefix

Rejecting/accepting selected routes
Set attributes to influence path selection

Tools:

Prefix-list (filter prefixes)

Filter-list (filter ASes)

Route-maps and communities

Policy Control Prefix List

- Filter routes based on prefix
- Inbound and Outbound

```
router bgp 200
neighbor 220.200.1.1 remote-as 210
neighbor 220.200.1.1 prefix-list PEER-IN in
neighbor 220.200.1.1 prefix-list PEER-OUT out
!
ip prefix-list PEER-IN deny 218.10.0.0/16
ip prefix-list PEER-IN permit 0.0.0.0/0 le 32
ip prefix-list PEER-OUT permit 215.7.0.0/16
```

Policy Control Filter List

- Filter routes based on AS path
- Inbound and Outbound

```
router bgp 100
neighbor 220.200.1.1 remote-as 210
neighbor 220.200.1.1 filter-list 5 out
neighbor 220.200.1.1 filter-list 6 in
!
ip as-path access-list 5 permit ^200$
ip as-path access-list 6 permit ^150$
```

Policy Control Regular Expressions

Like Unix regular expressions

- . Match one character
- * Match any number of preceding expression
- Match at least one of preceding expression
- ^ Beginning of line
- \$ End of line
- _ Beginning, end, white-space, brace
- Or
- () brackets to contain expression

Policy Control Regular Expressions

Simple Examples

.* Match anything

.+ Match at least one character

^\$ Match routes local to this AS

_1800\$ Originated by 1800

^1800_ Received from 1800

1800 Via 1800

_790_1800_ Passing through 1800 then 790

(1800)+ Match at least one of 1800 in sequence

\(65350\) Via 65350 (confederation AS)

- A route-map is like a "programme" for IOS
- Has "line" numbers, like programmes
- Each line is a separate condition/action
- Concept is basically:
 - if *match* then do *expression* and *exit* else
 - if *match* then do *expression* and *exit* else *etc*

Example using prefix-lists

```
router bgp 100
neighbor 1.1.1.1 route-map infilter in
route-map infilter permit 10
match ip address prefix-list HIGH-PREF
 set local-preference 120
route-map infilter permit 20
match ip address prefix-list LOW-PREF
set local-preference 80
route-map infilter permit 30
ip prefix-list HIGH-PREF permit 10.0.0.0/8
ip prefix-list LOW-PREF permit 20.0.0.0/8
```

Example using filter lists

```
router bgp 100
neighbor 220.200.1.2 route-map filter-on-as-path in
route-map filter-on-as-path permit 10
match as-path 1
 set local-preference 80
route-map filter-on-as-path permit 20
match as-path 2
set local-preference 200
route-map filter-on-as-path permit 30
ip as-path access-list 1 permit 150$
ip as-path access-list 2 permit 210
```

Example configuration of AS-PATH prepend

```
router bgp 300
network 215.7.0.0
neighbor 2.2.2.2 remote-as 100
neighbor 2.2.2.2 route-map SETPATH out
!
route-map SETPATH permit 10
set as-path prepend 300 300
```

 Use your own AS number when prepending Otherwise BGP loop detection may cause disconnects

Policy Control Setting Communities

Example Configuration

```
router bgp 100
neighbor 220.200.1.1 remote-as 200
neighbor 220.200.1.1 send-community
neighbor 220.200.1.1 route-map set-community out
route-map set-community permit 10
match ip address prefix-list NO-ANNOUNCE
 set community no-export
route-map set-community permit 20
ip prefix-list NO-ANNOUNCE permit 172.168.0.0/16 ge 17
```

Policy Control Matching Communities

Example Configuration

```
router bap 100
neighbor 220.200.1.2 remote-as 200
neighbor 220.200.1.2 route-map filter-on-community in
route-map filter-on-community permit 10
match community 1
 set local-preference 50
route-map filter-on-community permit 20
match community 2 exact-match
 set local-preference 200
ip community-list 1 permit 150:3 200:5
ip community-list 2 permit 88:6
```

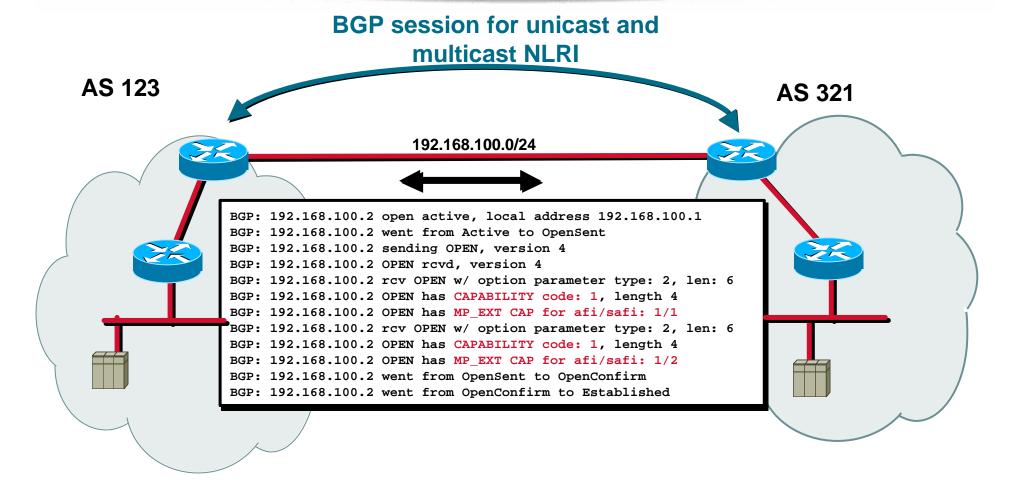


BGP Capabilities

- Documented in RFC2842
- Capabilities parameters passed in BGP open message
- Unknown or unsupported capabilities will result in NOTIFICATION message
- Current capabilities are:

0	Reserved	[RFC2842]
1	Multiprotocol Extensions for BGP-4	[RFC2858]
2	Route Refresh Capability for BGP-4	[RFC2918]
3	Cooperative Route Filtering Capability	[]
4	Multiple routes to a destination capability	[RFC3107]
64	Graceful Restart Capability	[]

BGP Capabilities Negotiation



BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

BGP Scaling Techniques CISCO SYSTEMS NANOG 22 © 2000, Cisco Systems, Inc.

BGP Scaling Techniques

- How to scale iBGP mesh beyond a few peers?
- How to implement new policy without causing flaps and route churning?
- How to reduce the overhead on the routers?
- How to keep the network stable, scalable, as well as simple?

BGP Scaling Techniques

- Dynamic Reconfiguration
- Peer groups
- Route flap damping
- Route Reflectors & Confederations

Dynamic Reconfiguration Soft Reconfiguration and Route Refresh www.cisco.com Presentation ID © 1999, Cisco Systems, Inc.

Soft Reconfiguration

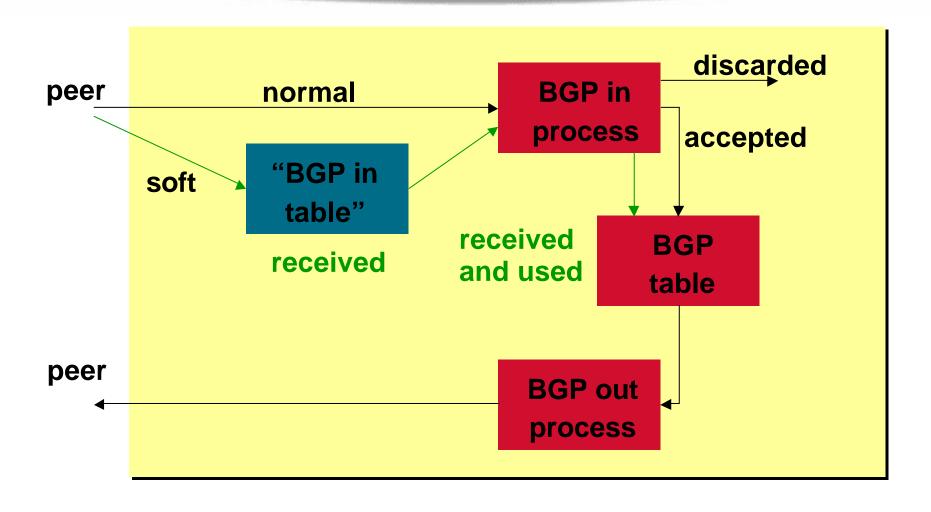
Problem:

- Hard BGP peer clear required after every policy change because the router does not store prefixes that are denied by a filter
- Hard BGP peer clearing consumes CPU and affects connectivity for all networks

Solution:

Soft-reconfiguration

Soft Reconfiguration



Soft Reconfiguration

- New policy is activated without tearing down and restarting the peering session
- Per-neighbour basis
- Use more memory to keep prefixes whose attributes have been changed or have not been accepted

Configuring Soft reconfiguration

```
router bgp 100
neighbor 1.1.1.1 remote-as 101
neighbor 1.1.1.1 route-map infilter in
neighbor 1.1.1.1 soft-reconfiguration inbound
```

! Outbound does not need to be configured!

Then when we change the policy, we issue an exec command

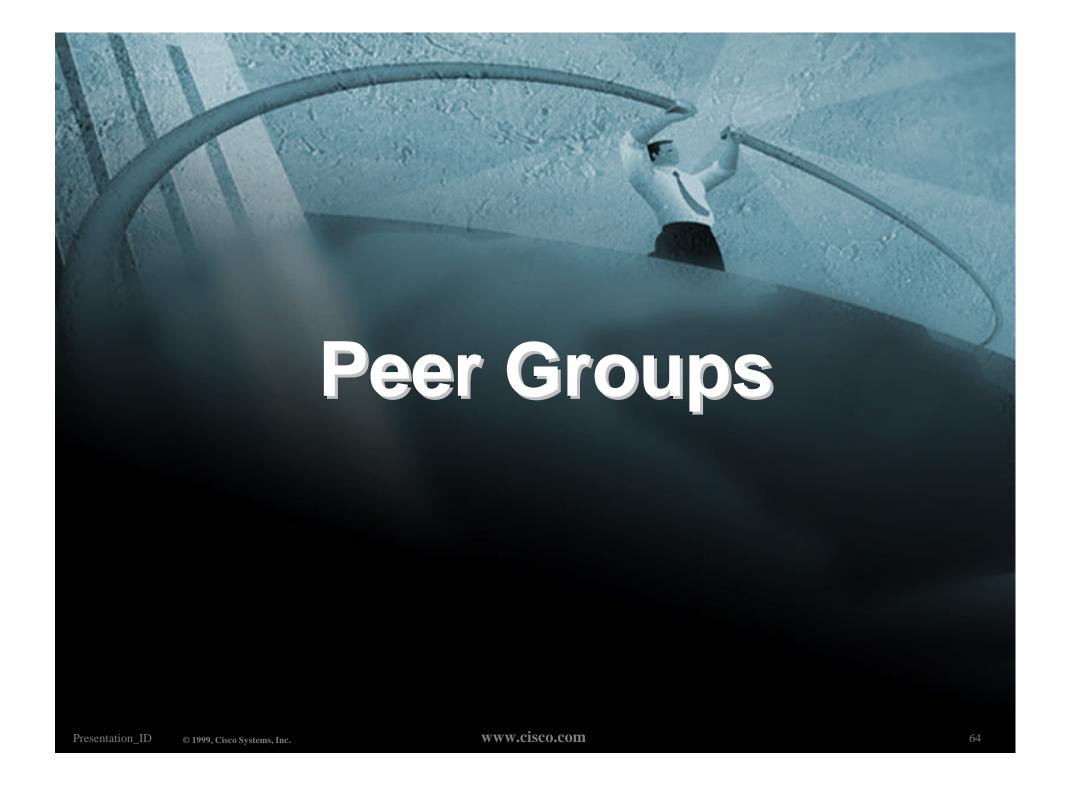
clear ip bgp 1.1.1.1 soft [in | out]

Route Refresh Capability

- Facilitates non-disruptive policy changes
- No configuration is needed
- No additional memory is used
- Requires peering routers to support "route refresh capability" – RFC2918
- clear ip bgp x.x.x.x in tells peer to resend full BGP announcement

Soft Reconfiguration vs Route Refresh

- Use Route Refresh capability if supported find out from "show ip bgp neighbor" uses much less memory
- Otherwise use Soft Reconfiguration
- Only hard-reset a BGP peering as a last resort



Peer Groups

Without peer groups

- iBGP neighbours receive same update
- Large iBGP mesh slow to build
- Router CPU wasted on repeat calculations
 Solution peer groups!
- Group peers with same outbound policy
- Updates are generated once per group

Peer Groups – Advantages

- Makes configuration easier
- Makes configuration less prone to error
- Makes configuration more readable
- Lower router CPU load
- iBGP mesh builds more quickly
- Members can have different inbound policy
- Can be used for eBGP neighbours too!

Configuring Peer Group

```
router bgp 100
 neighbor ibgp-peer peer-group
 neighbor ibgp-peer remote-as 100
 neighbor ibgp-peer update-source loopback 0
 neighbor ibgp-peer send-community
 neighbor ibgp-peer route-map outfilter out
 neighbor 1.1.1.1 peer-group ibgp-peer
 neighbor 2.2.2.2 peer-group ibgp-peer
 neighbor 2.2.2.2 route-map infilter in
 neighbor 3.3.3.3 peer-group ibgp-peer
! note how 2.2.2.2 has different inbound filter from peer-group!
```

Configuring Peer Group

```
router bap 109
neighbor external-peer peer-group
neighbor external-peer send-community
neighbor external-peer route-map set-metric out
neighbor 160.89.1.2 remote-as 200
neighbor 160.89.1.2 peer-group external-peer
neighbor 160.89.1.4 remote-as 300
neighbor 160.89.1.4 peer-group external-peer
neighbor 160.89.1.6 remote-as 400
neighbor 160.89.1.6 peer-group external-peer
neighbor 160.89.1.6 filter-list infilter in
```



Route Flap Damping

Route flap

Going up and down of path or change in attribute

BGP WITHDRAW followed by UPDATE = 1 flap eBGP neighbour going down/up is NOT a flap Ripples through the entire Internet

Wastes CPU

 Damping aims to reduce scope of route flap propagation

Route Flap Damping (Continued)

Requirements

Fast convergence for normal route changes

History predicts future behaviour

Suppress oscillating routes

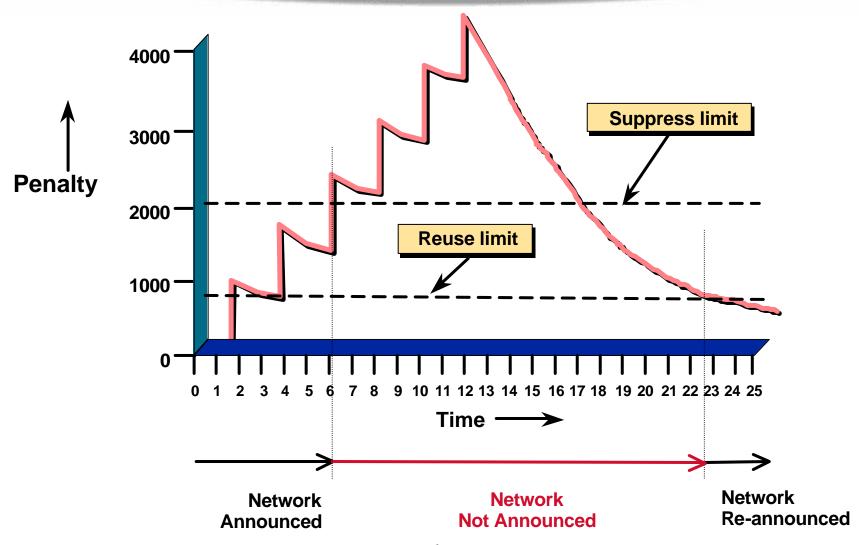
Advertise stable routes

Documented in RFC2439

Operation

- Add penalty (1000) for each flap
 Change in attribute gets penalty of 500
- Exponentially decay penalty half life determines decay rate
- Penalty above suppress-limit do not advertise route to BGP peers
- Penalty decayed below reuse-limit re-advertise route to BGP peers penalty reset to zero when it is half of reuse-limit

Operation



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Operation

- Only applied to inbound announcements from eBGP peers
- Alternate paths still usable
- Controlled by:

Half-life (default 15 minutes)

reuse-limit (default 750)

suppress-limit (default 2000)

maximum suppress time (default 60 minutes)

Configuration

Fixed damping

```
router bgp 100
```

```
bgp dampening [<half-life> <reuse-value> <suppress-
penalty> <maximum suppress time>]
```

Selective and variable damping

```
bgp dampening [route-map <name>]
```

Variable damping recommendations for ISPs

http://www.ripe.net/docs/ripe-210.html

Operation

- Care required when setting parameters
- Penalty must be less than reuse-limit at the maximum suppress time
- Maximum suppress time and half life must allow penalty to be larger than suppress limit

Configuration

Examples - *

bgp dampening 30 750 3000 60

reuse-limit of 750 means maximum possible penalty is 3000 – no prefixes suppressed as penalty cannot exceed suppress-limit

Examples - √

bgp dampening 30 2000 3000 60

reuse-limit of 2000 means maximum possible penalty is 8000 – suppress limit is easily reached

Maths!

Maximum value of penalty is

$$\frac{\left\langle \frac{\text{max-suppress-time}}{\text{half-life}} \right\rangle}{\text{max-penalty} = \text{reuse-limit} \times 2$$

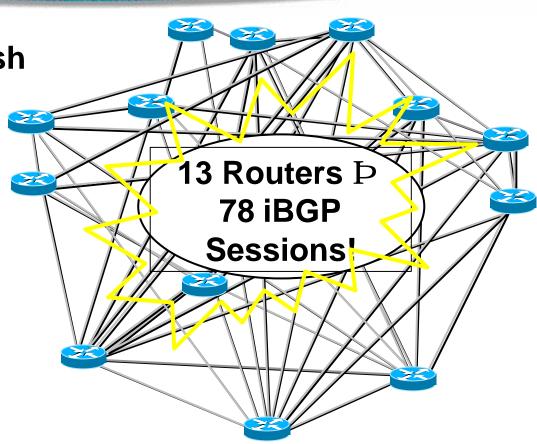
 Always make sure that suppress-limit is LESS than max-penalty otherwise there will be no flap damping



Scaling iBGP mesh

Avoid n(n-1)/2 iBGP mesh

n=1000 P nearly half a million ibgp sessions!

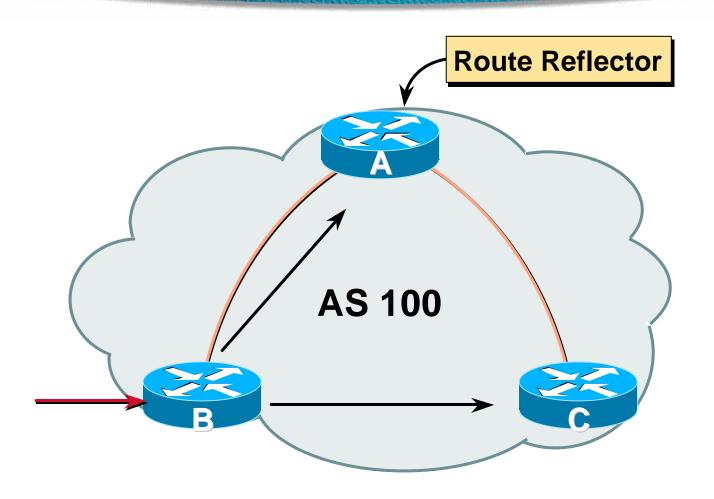


Two solutions

Route reflector – simpler to deploy and run

Confederation – more complex, corner case benefits

Route Reflector: Principle



Route Reflector

www.cisco.com

Reflector receives path from clients and non-clients

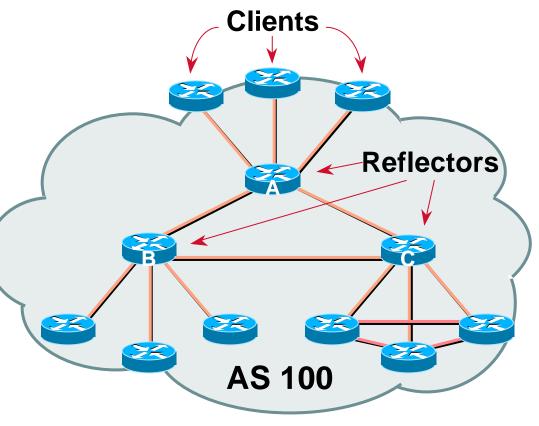
Selects best path

If best path is from client, reflect to other clients and non-clients

If best path is from non-client, reflect to clients only

Non-meshed clients

Described in RFC2796



Route Reflector Topology

- Divide the backbone into multiple clusters
- At least one route reflector and few clients per cluster
- Route reflectors are fully meshed
- Clients in a cluster could be fully meshed
- Single IGP to carry next hop and local routes

Route Reflectors: Loop Avoidance

Originator_ID attribute

Carries the RID of the originator of the route in the local AS (created by the RR)

Cluster_list attribute

The local cluster-id is added when the update is sent by the RR

Cluster-id is automatically set from routerid (address of loopback)

Do NOT use bgp cluster-id x.x.x.x

Route Reflectors: Redundancy

 Multiple RRs can be configured in the same cluster – not advised!

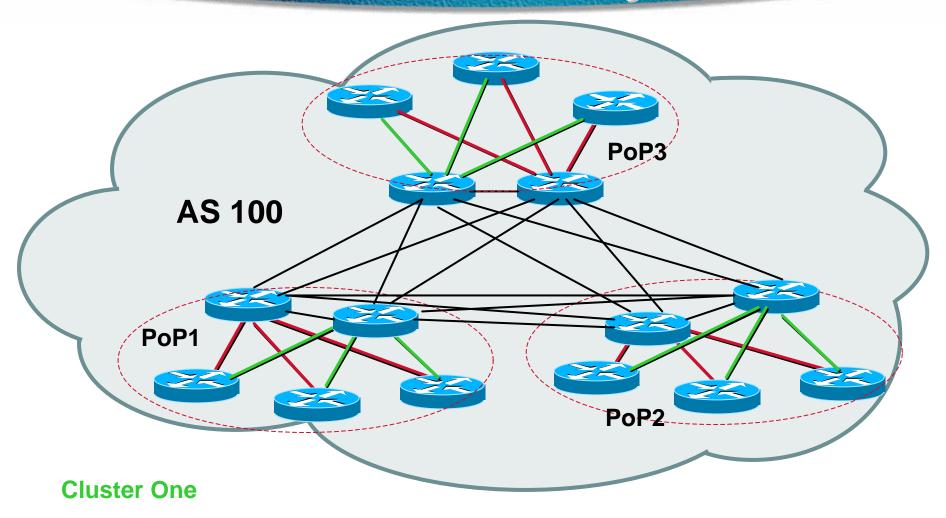
All RRs in the cluster must have the same cluster-id (otherwise it is a different cluster)

 A router may be a client of RRs in different clusters

Common today in ISP networks to overlay two clusters – redundancy achieved that way

® Each client has two RRs = redundancy

Route Reflectors: Redundancy



Cluster Two

Route Reflectors: Migration

• Where to place the route reflectors?

Always follow the physical topology!

This will guarantee that the packet forwarding won't be affected

Typical ISP network:

PoP has two core routers

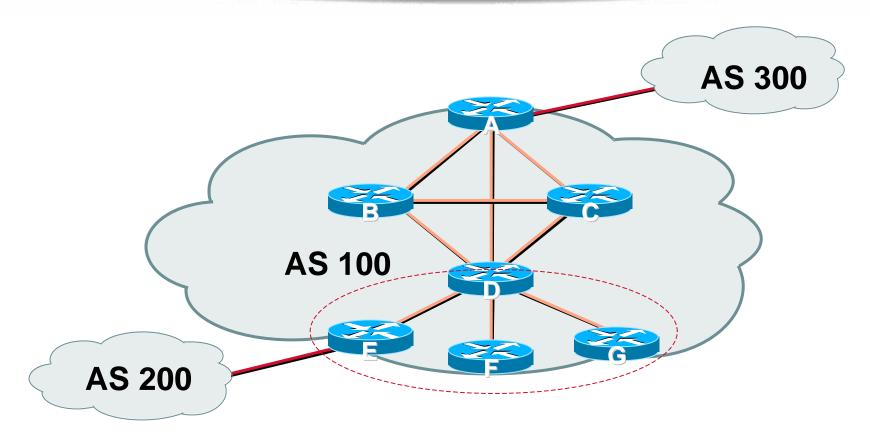
Core routers are RR for the PoP

Two overlaid clusters

Route Reflectors: Migration

- Typical ISP network:
 - Core routers have fully meshed iBGP
 Create further hierarchy if core mesh too big
 Split backbone into regions
- Configure one cluster pair at a time Eliminate redundant iBGP sessions Place maximum one RR per cluster Easy migration, multiple levels

Route Reflector: Migration



 Migrate small parts of the network, one part at a time.

Configuring a Route Reflector

```
neighbor 1.1.1.1 remote-as 100
neighbor 1.1.1.1 route-reflector-client
neighbor 2.2.2.2 remote-as 100
neighbor 2.2.2.2 route-reflector-client
neighbor 3.3.3.3 remote-as 100
neighbor 3.3.3.3 route-reflector-client
```

Confederations

Divide the AS into sub-AS

eBGP between sub-AS, but some iBGP information is kept

Preserve NEXT_HOP across the sub-AS (IGP carries this information)

Preserve LOCAL_PREF and MED

- Usually a single IGP
- Described in RFC3065

Confederations (Cont.)

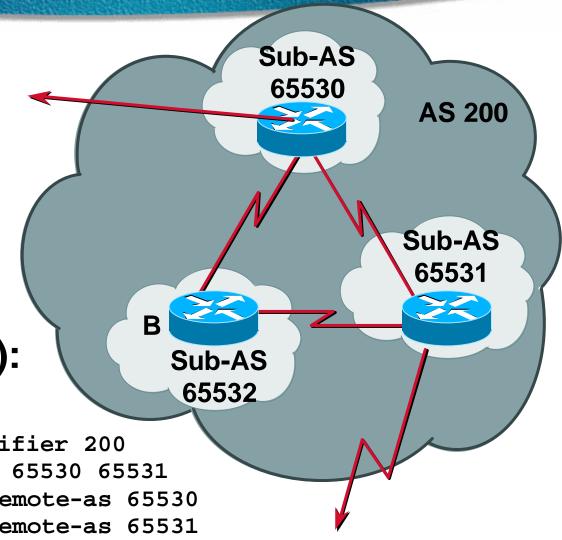
 Visible to outside world as single AS – "Confederation Identifier"

Each sub-AS uses a number from the private space (64512-65534)

 iBGP speakers in sub-AS are fully meshed

The total number of neighbors is reduced by limiting the full mesh requirement to only the peers in the sub-AS

Confederations (cont.)



• Configuration (rtr B):

router bgp 65532

bgp confederation identifier 200

bgp confederation peers 65530 65531

neighbor 141.153.12.1 remote-as 65530

neighbor 141.153.17.2 remote-as 65531

Route Propagation Decisions

• Same as with "normal" BGP:

From peer in same sub-AS \rightarrow only to external peers

From external peers \rightarrow to all neighbors

"External peers" refers to

Peers outside the confederation

Peers in a different sub-AS

Preserve LOCAL_PREF, MED and NEXT_HOP

Confederations (cont.)

• Example (cont.):

```
BGP table version is 78, local router ID is 141.153.17.1
Status codes: s suppressed, d damped, h history, * valid, >
best, i - internal
Origin codes: i - IGP, e - EGP, ? - incomplete
Network
                          Metric LocPrf Weight Path
              Next Hop
*> 10.0.0.0 141.153.14.3
                                 100
                                               (65531) 1 i
                            0
                                          0
*> 141.153.0.0 141.153.30.2
                                 100
                                               (65530) i
                            0
                                          0
*> 144.10.0.0 141.153.12.1
                                 100
                                          0
                                               (65530) i
                            0
*> 199.10.10.0 141.153.29.2
                                 100
                                          0
                                               (65530) 1 i
```

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RRs or Confederations

	Internet Connectivity	Multi-Level Hierarchy	Policy Control	Scalability	Migration Complexity
Confederations	Anywhere in the Network	Yes	Yes	Medium	Medium to High
Route Reflectors	Anywhere in the Network	Yes	Yes	Very High	Very Low

Most new service provider networks now deploy Route Reflectors from Day One

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More points about confederations

- Can ease "absorbing" other ISPs into you ISP – e.g., if one ISP buys another (can use local-as feature to do a similar thing)
- You can use route-reflectors with confederation sub-AS to reduce the sub-AS iBGP mesh

BGP Scaling Techniques

 These 4 techniques should be core requirements in all ISP networks

Soft reconfiguration/Route Refresh

Peer groups

Route flap damping

Route reflectors

BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

Deploying BGP in an ISP Network

Current Practices

CISCO SYSTEMS

BGP versus OSPF/ISIS

Internal Routing Protocols (IGPs)
 examples are ISIS and OSPF
 used for carrying infrastructure addresses

NOT used for carrying Internet prefixes or customer prefixes

design goal is to minimise number of prefixes in IGP to aid scalability and rapid convergence

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BGP versus OSPF/ISIS

- BGP used internally (iBGP) and externally (eBGP)
- iBGP used to carry some/all Internet prefixes across backbone customer prefixes
- eBGP used to exchange prefixes with other ASes implement routing policy

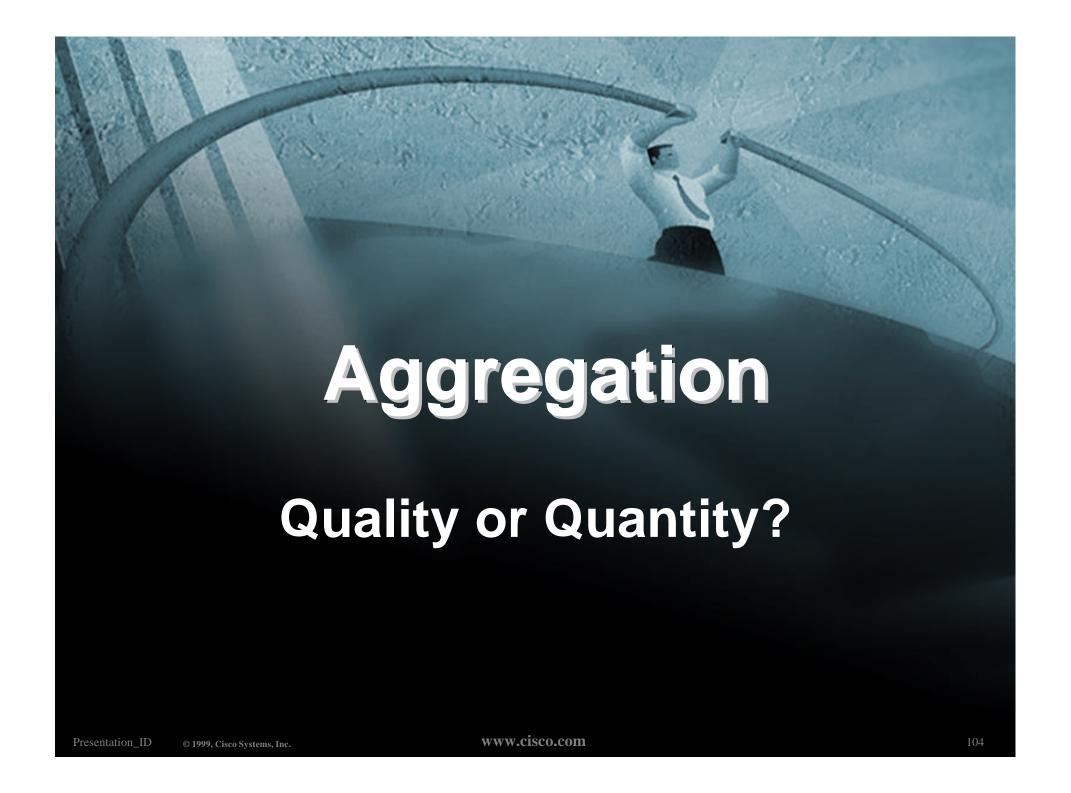
BGP versus OSPF/ISIS

DO NOT:

distribute BGP prefixes into an IGP distribute IGP routes into BGP use an IGP to carry customer prefixes

YOUR NETWORK WILL NOT SCALE

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Aggregation

- ISPs receive address block from Regional Registry or upstream provider
- Aggregation means announcing the address block only, not subprefixes
 - Subprefixes should only be announced in special cases see later.
- Aggregate should be generated internally Not on the network borders!

Configuring Aggregation – Method One

- ISP has 221.10.0.0/19 address block
- To put into BGP as an aggregate:

```
router bgp 100
network 221.10.0.0 mask 255.255.224.0
ip route 221.10.0.0 255.255.224.0 null0
```

 The static route is a "pull up" route more specific prefixes within this address block ensure connectivity to ISP's customers

"longest match lookup"

Configuring Aggregation – Method Two

Configuration Example

```
router bgp 109
network 221.10.0.0 mask 255.255.252.0
aggregate-address 221.10.0.0 255.255.224.0 [summary-only]
```

- Requires more specific prefix in routing table before aggregate is announced
- {summary-only} keyword
 - ensures that only the summary is announced if a more specific prefix exists in the routing table
- Sets "aggregator" attribute
 Useful for debugging

Announcing Aggregate – Cisco IOS

Configuration Example

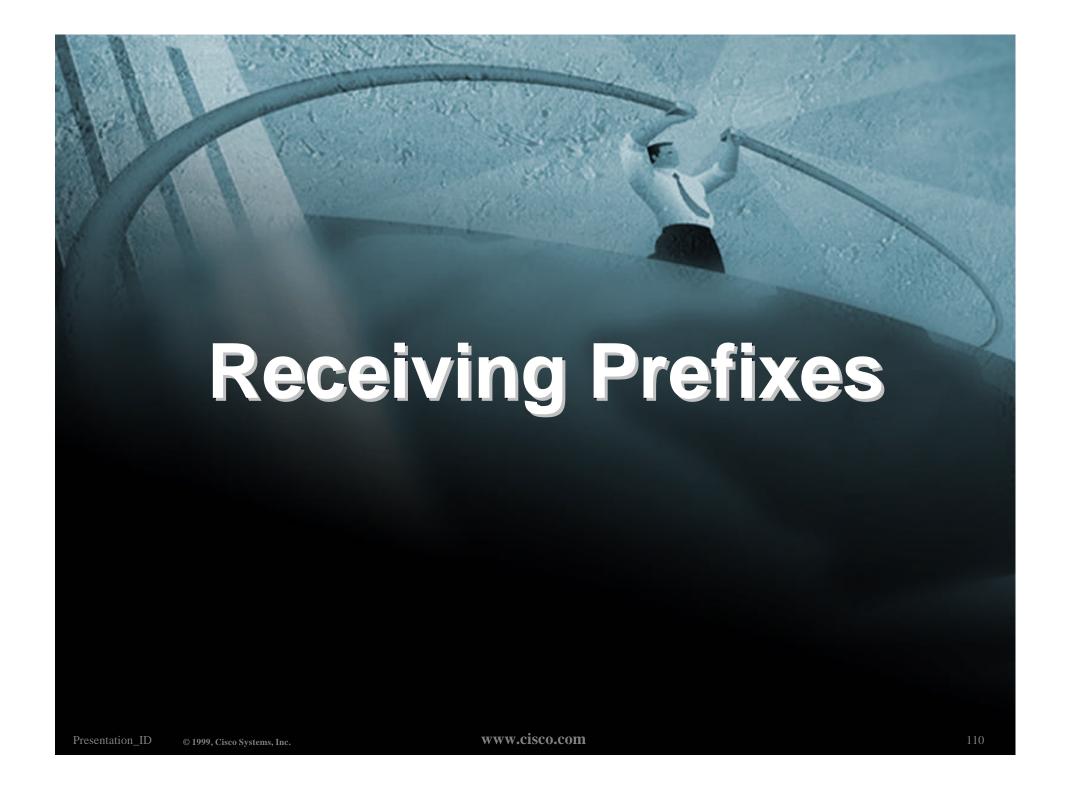
```
router bgp 100
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 101
neighbor 222.222.10.1 prefix-list out-filter out
!
ip route 221.10.0.0 255.255.224.0 null0
!
ip prefix-list out-filter permit 221.10.0.0/19
```

Announcing an Aggregate

- ISPs who don't and won't aggregate are held in poor regard by community
- Registries' minimum allocation size is now a /20

no real reason to see subprefixes of allocated blocks in the Internet

BUT there are currently >60000 /24s!



- ISPs should only accept prefixes which have been assigned or allocated to their downstream peer
- For example
 downstream has 220.50.0.0/20 block
 should only announce this to peers
 peers should only accept this from them

Receiving Prefixes – Cisco IOS

Configuration Example on upstream

```
router bgp 100
neighbor 222.222.10.1 remote-as 101
neighbor 222.222.10.1 prefix-list customer in
!
ip prefix-list customer permit 220.50.0.0/20
```

- Not desirable unless really necessary special circumstances – see later
- Ask upstream to either:
 originate a default-route
 announce one prefix you can use as default

Downstream Router Configuration

```
router bgp 100
network 221.10.0.0 mask 255.255.224.0
neighbor 221.5.7.1 remote-as 101
neighbor 221.5.7.1 prefix-list infilter in
neighbor 221.5.7.1 prefix-list outfilter out
!
ip prefix-list infilter permit 0.0.0.0/0
!
ip prefix-list outfilter permit 221.10.0.0/19
```

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Upstream Router Configuration

```
router bgp 101
neighbor 221.5.7.2 remote-as 100
neighbor 221.5.7.2 default-originate
neighbor 221.5.7.2 prefix-list cust-in in
neighbor 221.5.7.2 prefix-list cust-out out
!
ip prefix-list cust-in permit 221.10.0.0/19
!
ip prefix-list cust-out permit 0.0.0.0/0
```

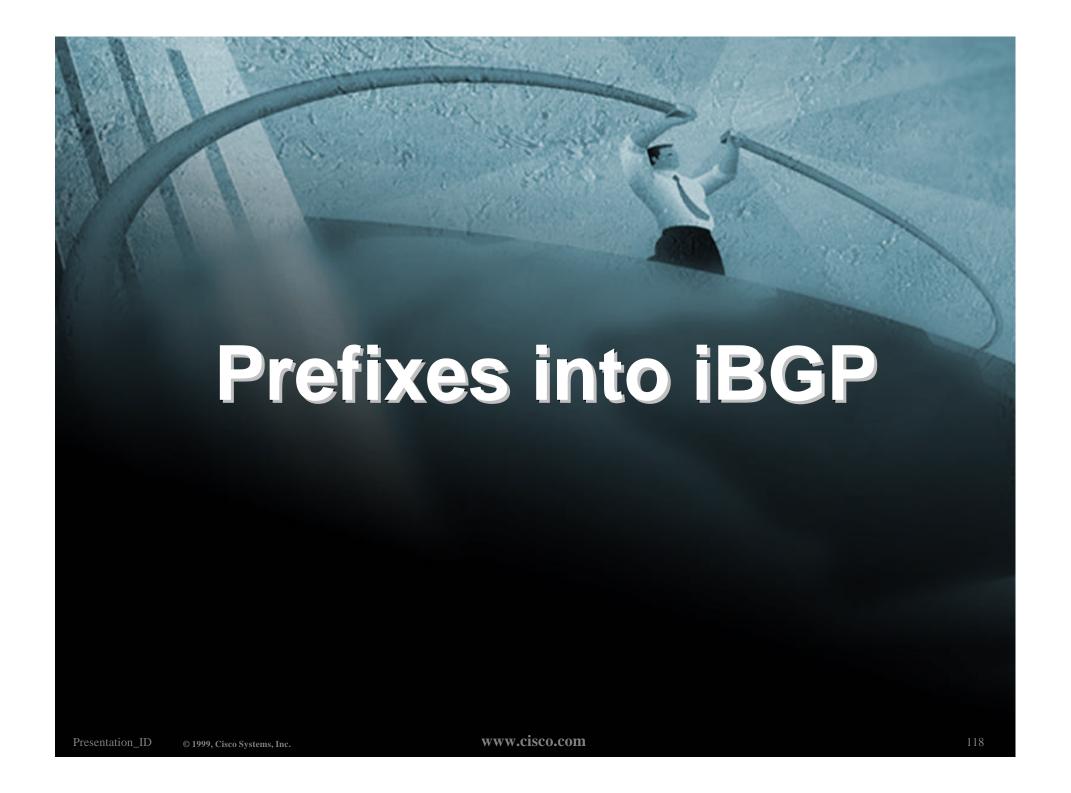
 If necessary to receive prefixes from upstream provider, care is required don't accept RFC1918 etc prefixes

http://www.ietf.org/internet-drafts/draft-manning-dsua-06.txt

don't accept your own prefix don't accept default (unless you need it) don't accept prefixes longer than /24

Receiving Prefixes

```
router bgp 100
 network 221.10.0.0 mask 255.255.224.0
neighbor 221.5.7.1 remote-as 101
 neighbor 221.5.7.1 prefix-list in-filter in
ip prefix-list in-filter deny 0.0.0.0/0
                                            ! Block default
ip prefix-list in-filter deny 0.0.0.0/8 le 32
ip prefix-list in-filter deny 10.0.0.0/8 le 32
ip prefix-list in-filter deny 127.0.0.0/8 le 32
ip prefix-list in-filter deny 169.254.0.0/16 le 32
ip prefix-list in-filter deny 172.16.0.0/12 le 32
ip prefix-list in-filter deny 192.0.2.0/24 le 32
ip prefix-list in-filter deny 192.168.0.0/16 le 32
ip prefix-list in-filter deny 221.10.0.0/19 le 32 ! Block local prefix
ip prefix-list in-filter deny 224.0.0.0/3 le 32 ! Block multicast
ip prefix-list in-filter deny 0.0.0.0/0 ge 25
                                                 ! Block prefixes >/24
ip prefix-list in-filter permit 0.0.0.0/0 le 32
```



Injecting prefixes into iBGP

- Use iBGP to carry customer prefixes don't ever use IGP
- Point static route to customer interface
- Use BGP network statement
- As long as static route exists (interface active), prefix will be in BGP

Router Configuration network statement

• Example:

```
interface loopback 0
 ip address 215.17.3.1 255.255.255.255
interface Serial 5/0
 ip unnumbered loopback 0
 ip verify unicast reverse-path
ip route 215.34.10.0 255.255.252.0 Serial 5/0
router bgp 100
network 215.34.10.0 mask 255.255.252.0
```

Injecting prefixes into iBGP

 interface flap will result in prefix withdraw and re-announce

use "ip route...permanent"

Static route always exists, even if interface is down ® prefix announced in iBGP

 many ISPs use redistribute static rather than network statement

only use this if you understand why

Inserting prefixes into BGP – redistribute static

Care required with redistribute!

redistribute <routing-protocol> means everything in the <routing-protocol> will be transferred into the current routing protocol

Does not scale if uncontrolled

Best avoided if at all possible

redistribute normally used with "route-maps" and under tight administrative control

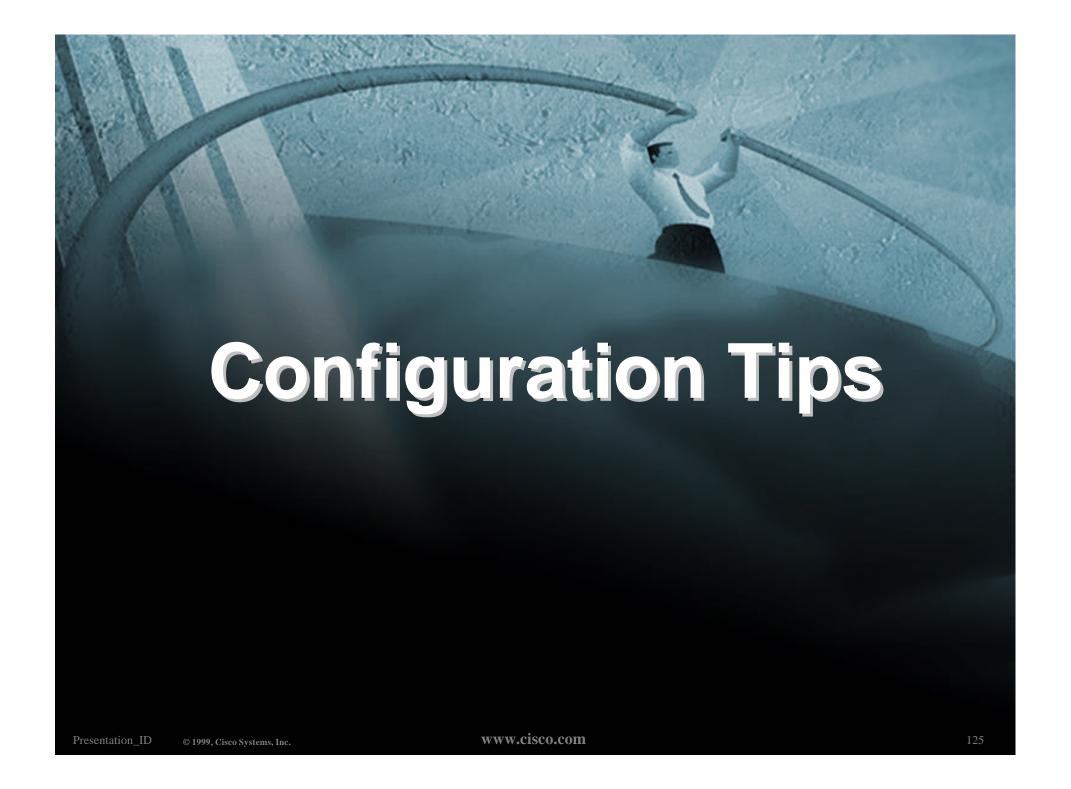
Router Configuration redistribute static

• Example:

```
ip route 215.34.10.0 255.255.252.0 Serial 5/0
router bgp 100
 redistribute static route-map static-to-bgp
<snip>
route-map static-to-bgp permit 10
match ip address prefix-list ISP-block
 set origin igp
<snip>
ip prefix-list ISP-block permit 215.34.10.0/22 le 30
```

Injecting prefixes into iBGP

- Route-map ISP-block can be used for many things:
 - setting communities and other attributes setting origin code to IGP, etc
- Be careful with prefix-lists and route-maps absence of either/both could mean all statically routed prefixes go into iBGP



iBGP and IGPs

Make sure loopback is configured on router

iBGP between loopbacks, NOT real interfaces

- Make sure IGP carries loopback /32 address
- Make sure IGP carries DMZ nets
 Or use next-hop-self on iBGP neighbours

neighbor x.x.x.x next-hop-self

Next-hop-self

Used by many ISPs on edge routers

Preferable to carrying DMZ /30 addresses in the IGP

Reduces size of IGP to just core infrastructure

Alternative to using ip unnumbered

Helps scale network

BGP speaker announces external network using local address (loopback) as next-hop

BGP Template – iBGP peers

iBGP Peer Group

AS100 router bgp 100 neighbor internal peer-group neighbor internal description ibgp peers neighbor internal remote-as 100 neighbor internal update-source Loopback0 neighbor internal next-hop-self neighbor internal send-community neighbor internal version 4 neighbor internal password 7 03085A09 neighbor 1.0.0.1 peer-group internal neighbor 1.0.0.2 peer-group internal

BGP Template – iBGP peers

- Use peer-groups
- iBGP between loopbacks!
- Next-hop-self
 Keep DMZ and point-to-point out of IGP
- Always send communities in iBGP
 Otherwise accidents will happen
- Hardwire BGP to version 4
 Yes, this is being paranoid!
- Use passwords on iBGP session
 Not being paranoid, VERY necessary

BGP Template – eBGP peers

Router B: AS 200 router bgp 100 10.0.0.0 bgp dampening route-map RIPE-210-flap network 10.60.0.0 mask 255.255.0.0 neighbor external peer-group **AS 100 is a** neighbor external remote-as 200 customer neighbor external description ISP connection of AS 200 neighbor external remove-private-AS neighbor external version 4 10.200.0.0 neighbor external prefix-list is pout out; "accident" filter neighbor external route-map ispout out; "real" filter neighbor external route-map ispin in 10.60.0.0/16 neighbor external password 7 020A0559 **AS100** neighbor external maximum-prefix 120000 [warning-only] neighbor 10.200.0.1 peer-group external ip route 10.60.0.0 255.255.0.0 null0 254

В

BGP Template – eBGP peers

- BGP damping use RIPE-210 parameters
- Remove private ASes from announcements
 Common omission today
- Use extensive filters, with "backup"
- Use password agreed between you and peer on eBGP session
- Use maximum-prefix tracking

Router will warn you if there are sudden changes in BGP table size, bringing down eBGP if necessary

More BGP "defaults"

Log neighbour changes

bgp log-neighbor-changes

Enable deterministic MED

bgp deterministic-med

Otherwise bestpath could be different every time BGP session is reset

Make BGP admin distance higher than any IGP

distance bgp 200 200 200

Customer Aggregation

BGP customers

Offer max 3 types of feeds (easier than custom configuration per peer)

Use communities

Static customers

Use communities

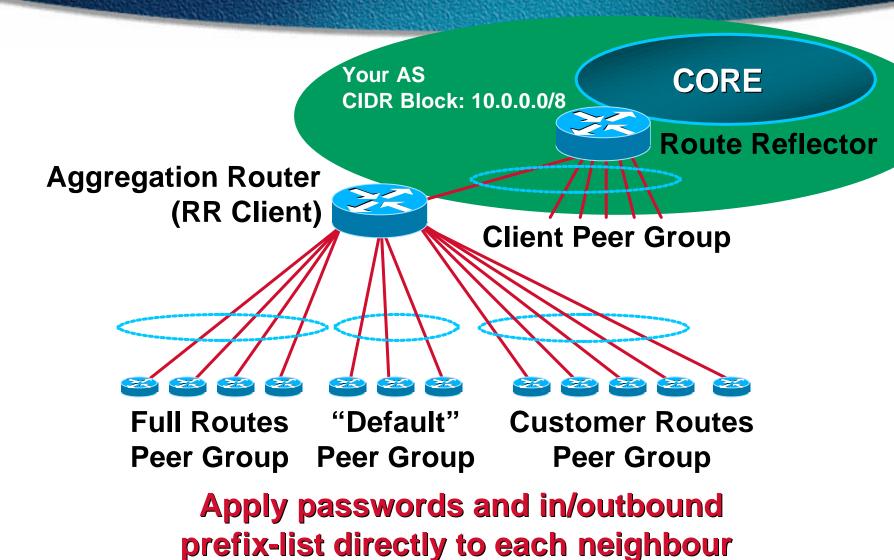
Differentiate between different types of prefixes

Makes eBGP filtering easy

BGP Customer Aggregation Guidelines

- Define at least three peer groups: cust-default—send default route only cust-cust—send customer routes only cust-full —send full Internet routes
- Identify routes via communities e.g.
 100:4100=customers; 100:4500=peers
- Apply passwords per neighbour
- Apply inbound & outbound prefix-list per neighbour

BGP Customer Aggregation



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Static Customer Aggregation Guidelines

Identify routes via communities, e.g.

100:4000=my address blocks

100:4200=customers from my block

100:4300=customers outside my block

Helps with aggregation, iBGP, filtering

 BGP network statements on aggregation routers set correct community

Sample core configuration

eBGP peers and upstreams

Send communities 100:4000, 100:4100 and 100:4300, receive everything

iBGP full routes

Send everything (only network core)

iBGP partial routes

Send communities 100:4000, 100:4100, 100:4200, 100:4300 and 100:4500 (edge routers, peering routers, IXP routers)

Simple configuration with peer-groups and route-maps

Acquisitions!

- Your ISP has just bought another ISP How to merge networks?
- Options:

use confederations – make their AS a sub-AS (only useful if you are using confederations already)

use the BGP local-as feature to implement a gradual transition – overrides BGP process ID

neighbor x.x.x.x local-as as-number

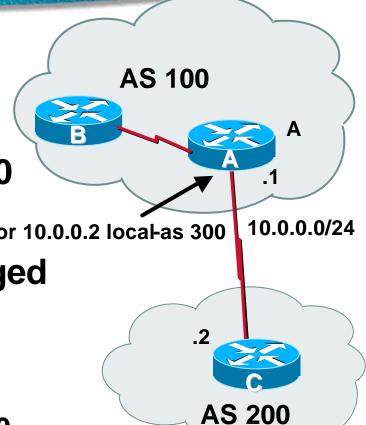
local-AS - Application

Router A has a process ID of 100

• The peering with AS200 is neighbor 10.0.0.2 local-as 300 established as if router A belonged to AS300.

AS_PATH

routes originated in AS100 = 300 100 routes received from AS200 = 300 200



BGP for Internet Service Providers

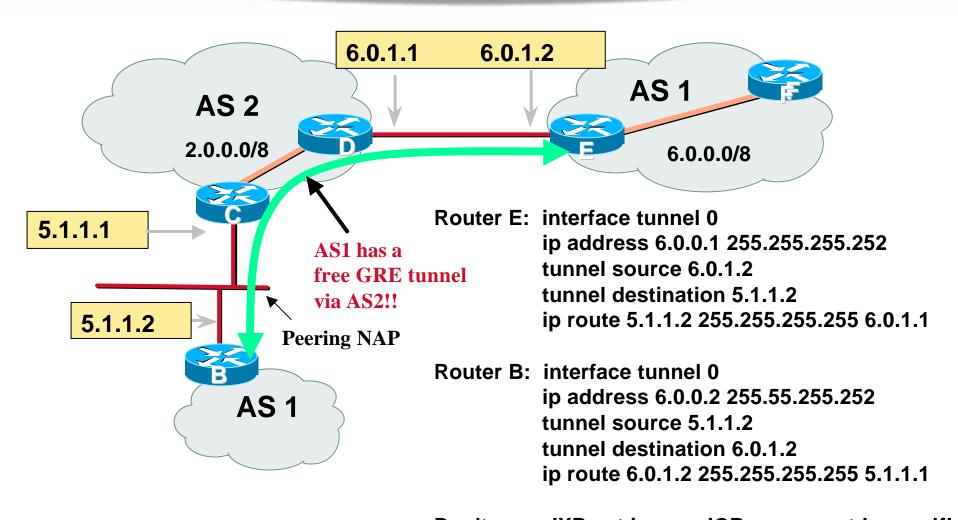
- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

Troubleshooting Staying out of Trouble CISCO SYSTEMS NANOG 22 © 2000, Cisco Systems, Inc.

Potential Caveats and Operational Problems

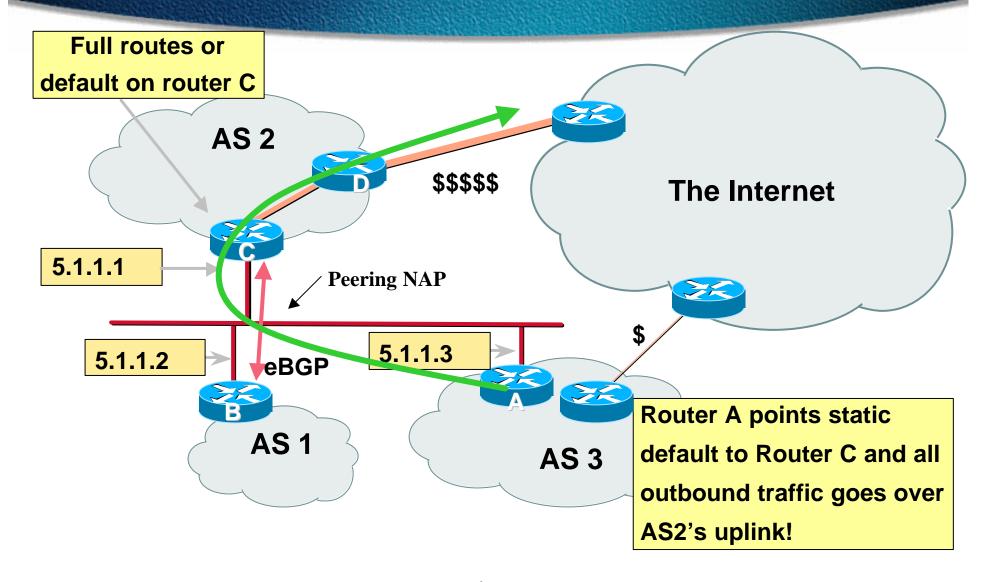
- GRE Tunnels & IXPs
- Auto-summarisation & synchronisation
- Route Reflectors
 Follow the topology
- Common Problems
 ...and the solutions!

Prevent GRE VPNs



Don't carry IXP net in your IGP – use next-hop-self!

Prevent "Defaulting"



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Watch out at IXPs/NAPs

- IXP router should not carry full routes or have a default
- ISP should not carry IXP/NAP network prefix internally

Use BGP next-hop-self

- or -
- Use RPF check for non-peers
- Use good filters for peers

Auto Summarisation - Cisco IOS

- Historical feature
- Automatically summarises subprefixes to the classful network for prefixes redistributed into BGP

Example:

```
61.10.8.0/22 --> 61.0.0.0/8
```

 Must be turned off for any Internet connected site using BGP.

```
router bgp 109 no auto-summary
```

Synchronisation - Cisco IOS

- Historical feature
- BGP will not advertise a route before all routers in the AS have learned it via an IGP
- Disable synchronisation if:

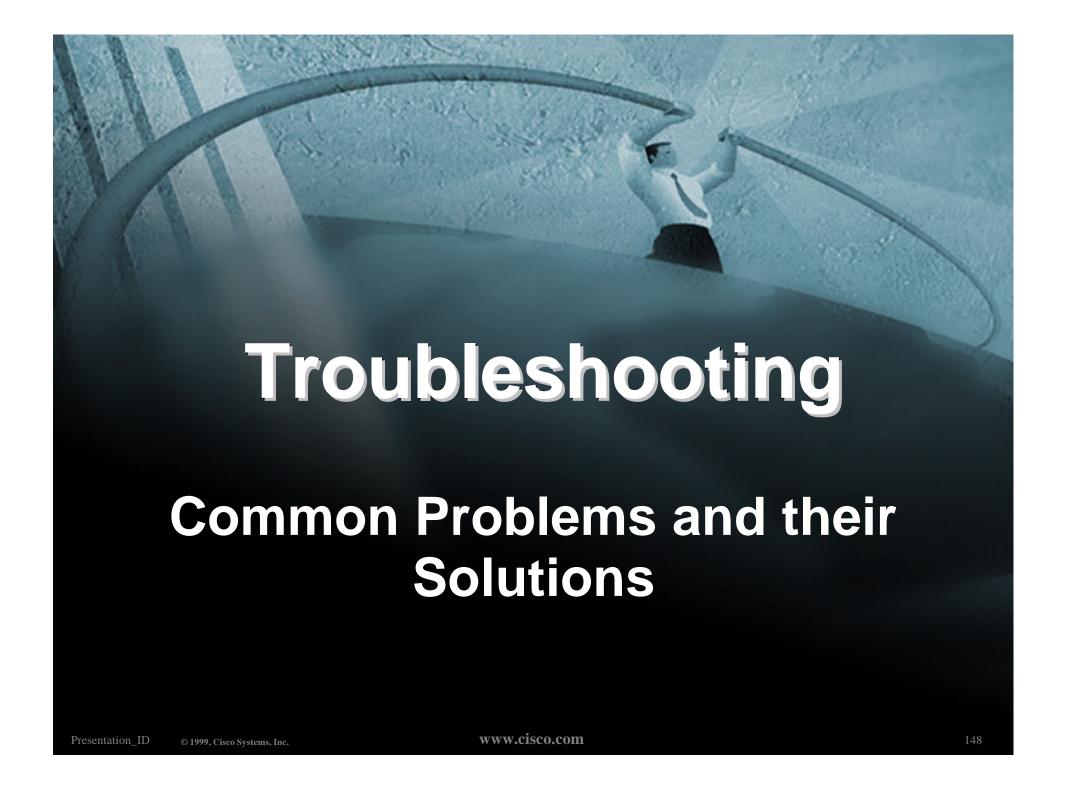
AS doesn't pass traffic from one AS to another, or

All transit routers in AS run BGP, or

iBGP is used across backbone

router bgp 109

no synchronization



Troubleshooting – Examples

- Missing routes
- Route Oscillation
- Routing Loops
- Troubleshooting hints

Route Origination

Network statement with mask

```
R1# show run | begin bgp
network 200.200.0.0 mask 255.255.252.0
```

BGP is not originating the route???

```
R1# show ip bgp | include 200.200.0.0
R1#
```

Do we have the exact route?

```
R1# show ip route 200.200.0.0 255.255.252.0 % Network not in table
```

Route Origination

Nail down routes you want to originate

```
R1#ip route 200.200.0.0 255.255.252.0 Null 0 200
```

Check the RIB

```
R1# show ip route 200.200.0.0 255.255.252.0

200.200.0.0/22 is subnetted, 1 subnets

S 200.200.0.0 [1/0] via Null 0
```

BGP originates the route!!

Route Oscillation

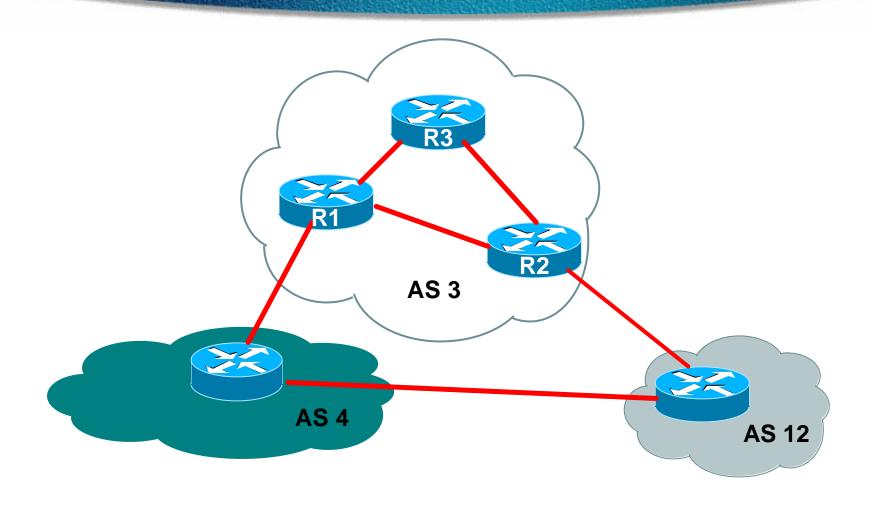
• One of the most common problems!

Every minute routes flap in the routing table from one next hop to another

With large routing table the most obvious symptom is high CPU in the "BGP-Router" process

Can be frustrating to track down unless you have seen it before!

Route Oscillation - Diagram



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Route Oscillation - Symptom

```
R3#show ip bgp summary
BGP router identifier 3.3.3.3, local AS number 3
BGP table version is 502, main routing table version 502
267 network entries and 272 paths using 34623 bytes of memory
R3#sh ip route summary | begin bgp
bap 3
                             6
                                             520
                                                         1400
  External: 0 Internal: 10 Local: 0
internal
                                                              5800
                                         13936
Total
                             263
                                                     43320
                10
```

Watch for:

table version number incrementing rapidly number of networks/paths or external/internal routes changing.

Pick up a bgp route from the RIB that is less than a minute old and watch what happens with the routing/bgp table ...

```
R3#show ip route 156.1.0.0
Routing entry for 156.1.0.0/16
  Known via "bgp 3", distance 200, metric 0
 Routing Descriptor Blocks:
  * 1.1.1.1, from 1.1.1.1, 00:00:53 ago
      Route metric is 0, traffic share count is 1
      AS Hops 2, BGP network version 474
R3#show ip bgp 156.1.0.0
BGP routing table entry for 156.1.0.0/16, version 474
Paths: (2 available, best #1)
  Advertised to non peer-group peers:
    2.2.2.2
  4 12
    1.1.1.1 from 1.1.1.1 (1.1.1.1)
      Origin IGP, localpref 100, valid, internal, best
  12
    142.108.10.2 (inaccessible) from 2.2.2.2 (2.2.2.2)
      Origin IGP, metric 0, localpref 100, valid, internal
```

...and after bgp_scanner runs (by default once a minute):

```
R3#sh ip route 156.1.0.0
Routing entry for 156.1.0.0/16
  Known via "bgp 3", distance 200, metric 0
    Routing Descriptor Blocks:
  * 142.108.10.2, from 2.2.2.2, 00:00:27 ago
      Route metric is 0, traffic share count is 1
      AS Hops 1, BGP network version 478
R3#sh ip bgp 156.1.0.0
BGP routing table entry for 156.1.0.0/16, version 478
Paths: (2 available, best #2)
  Advertised to non peer-group peers:
    1.1.1.1
  4 12
    1.1.1.1 from 1.1.1.1 (1.1.1.1)
      Origin IGP, localpref 100, valid, internal
  12
    142.108.10.2 from 2.2.2.2 (2.2.2.2)
      Origin IGP, metric 0, localpref 100, valid, internal, best
```

Let's take a look at the next hop at this point!

```
R3#show ip route 142.108.10.2
Routing entry for 142.108.0.0/16
  Known via "bgp 3", distance 200, metric 0
 Routing Descriptor Blocks:
  * 142.108.10.2, from 2.2.2.2, 00:00:50 ago
      Route metric is 0, traffic share count is 1
      AS Hops 1, BGP network version 476
R3#show ip bgp 142.108.10.2
BGP routing table entry for 142.108.0.0/16, version 476
Paths: (2 available, best #2)
  Advertised to non peer-group peers:
    1.1.1.1
  4 12
    1.1.1.1 from 1.1.1.1 (1.1.1.1)
      Origin IGP, localpref 100, valid, internal
  12
    142.108.10.2 from 2.2.2.2 (2.2.2.2)
      Origin IGP, metric 0, localpref 100, valid, internal, best
```

Next-hop is recursive !!! This will be detected next time the scanner runs and the other path will be installed in the RIB instead

```
R3#sh debug

BGP events debugging is on

BGP updates debugging is on

IP routing debugging is on

R3#

BGP: scanning routing tables

BGP: nettable_walker 142.108.0.0/16 calling revise_route

RT: del 142.108.0.0 via 142.108.10.2, bgp metric [200/0]

BGP: revise route installing 142.108.0.0/16 -> 1.1.1.1

RT: add 142.108.0.0/16 via 1.1.1.1, bgp metric [200/0]

RT: del 156.1.0.0 via 142.108.10.2, bgp metric [200/0]

BGP: revise route installing 156.1.0.0/16 -> 1.1.1.1

RT: add 156.1.0.0/16 via 1.1.1.1, bgp metric [200/0]
```

The route to the next-hop is now valid and at the next bgp scan we will change to the shorter as-path path, and so on ...

```
R3#
BGP: scanning routing tables
BGP: ip nettable_walker 142.108.0.0/16 calling revise_route
RT: del 142.108.0.0 via 1.1.1.1, bgp metric [200/0]
BGP: revise route installing 142.108.0.0/16 -> 142.108.10.2
RT: add 142.108.0.0/16 via 142.108.10.2, bgp metric [200/0]
BGP: nettable_walker 156.1.0.0/16 calling revise_route
RT: del 156.1.0.0 via 1.1.1.1, bgp metric [200/0]
BGP: revise route installing 156.1.0.0/16 -> 142.108.10.2
RT: add 156.1.0.0/16 via 142.108.10.2, bgp metric [200/0]
```

Route Oscillation - Summary

- iBGP preserves the next-hop information from eBGP
- To avoid problems
 use "next-hop-self" for iBGP peering

-or-

make sure you advertise the next-hop prefix via the IGP

Inconsistent Route Selection

- Two common problems with route selection Inconsistency
 Appearance of an Incorrect decision
- RFC 1771 defines the decision algorithm
- Every vendor has tweaked the algorithm http://www.cisco.com/warp/public/459/25.shtml
- Route Selection problems can result from oversights in RFC1771

Inconsistent Route Selection

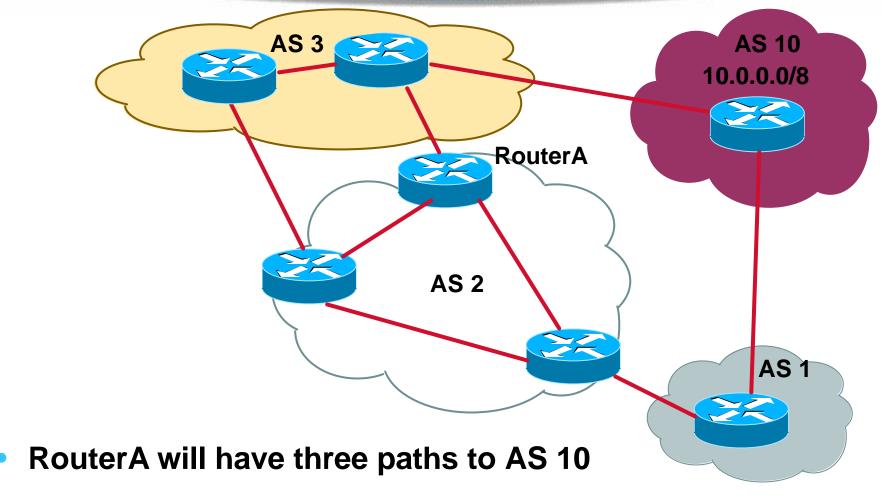
- RFC says that MED is not always compared
- As a result, the ordering of the paths can affect the decision process
- By default, the prefixes are compared in order of arrival (most recent to oldest)

use bgp deterministic-med to order paths consistently

the bestpath is recalculated as soon as the command is entered

enable in all the routers in the AS

Symptom - Diagram



MEDs from AS 3 will not be compared with MEDs from AS 1

Inconsistent Route Selection

```
RouterA#sh ip bgp 10.0.0.0

BGP routing table entry for 10.0.0.0/8, version 40

Paths: (3 available, best #3, advertised over IBGP, EBGP)

3 10
2.2.2.2 from 2.2.2.2
Origin IGP, metric 20, localpref 100, valid, internal

3 10
3.3.3.3 from 3.3.3.3
Origin IGP, metric 30, valid, external

1 10
1.1.1.1 from 1.1.1.1
Origin IGP, metric 0, localpref 100, valid, internal, best
```

Initial State

Path 1 beats Path 2 – Lower MED

Path 3 beats Path 1 – Lower Router-ID

Inconsistent Route Selection

```
RouterA#sh ip bgp 10.0.0.0

BGP routing table entry for 10.0.0.0/8, version 40

Paths: (3 available, best #3, advertised over IBGP, EBGP)

1 10
1.1.1.1 from 1.1.1.1
Origin IGP, metric 0, localpref 100, valid, internal

3 10
2.2.2.2 from 2.2.2.2
Origin IGP, metric 20, localpref 100, valid, internal

3 10
3.3.3.3 from 3.3.3.3
Origin IGP, metric 30, valid, external, best
```

1.1.1.1 bounced so the paths are re-ordered
 Path 1 beats Path 2 – Lower Router-ID
 Path 3 beats Path 1 – External vs Internal

Deterministic MED – Operation

- The paths are ordered by Neighbour AS
- The bestpath for each Neighbour AS group is selected
- The overall bestpath results from comparing the winners from each group
- The bestpath will be consistent because paths will be placed in a deterministic order

Deterministic MED – Result

```
RouterA#sh ip bgp 10.0.0.0

BGP routing table entry for 10.0.0.0/8, version 40

Paths: (3 available, best #1, advertised over IBGP, EBGP)

1 10
1.1.1.1 from 1.1.1.1
Origin IGP, metric 0, localpref 100, valid, internal, best

3 10
2.2.2.2 from 2.2.2.2
Origin IGP, metric 20, localpref 100, valid, internal

3 10
3.3.3.3 from 3.3.3.3
Origin IGP, metric 30, valid, external
```

Path 1 is best for AS 1

Path 2 beats Path 3 for AS 3 – Lower MED

Path 1 beats Path 2 – Lower Router-ID

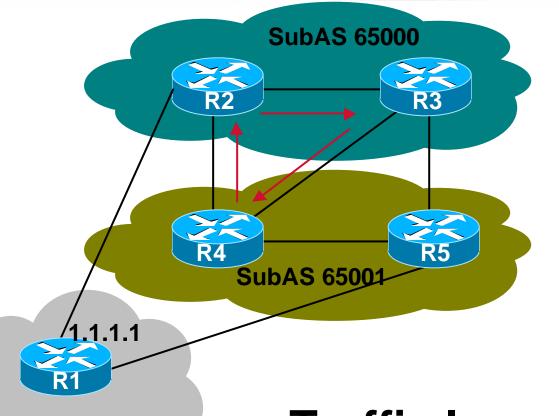
Deterministic MED – Summary

 If multihoming with multiple ISPs and peering with one ISP at multiple points:

use "bgp deterministic-med" enable it on all routers in the AS

Always use "bgp deterministic-med"

Routing Loop - Problem



10.0.0.0/8

SubAS 65002

traceroute 10.1.1.1

1 30,100,1,1

2 20.20.20.4 - R3

3 30.1.1.26 - R4

4 30.1.1.17 - R2

5 20.20.20.4 - R3

6 30.1.1.26 - R4

7 30.1.1.17 - R2

8 20.20.20.4

9 30.1.1.26

10 30.1.1.17

 Traffic loops between R3, R4, and R2

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Routing Loop – Diagnosis

- First grab a "show ip route" from the three problem routers
- R3 is forwarding traffic to 1.1.1.1 (R1)

```
R3# show ip route 10.1.1.1

Routing entry for 10.0.0.0/8

Known via "bgp 65000", distance 200, metric 0

Routing Descriptor Blocks:

1.1.1.1, from 5.5.5.5, 01:46:43 ago

Route metric is 0, traffic share count is 1

AS Hops 0, BGP network version 0

* 1.1.1.1, from 4.4.4.4, 01:46:43 ago

Route metric is 0, traffic share count is 1

AS Hops 0, BGP network version 0
```

Routing Loop – Diagnosis

R4 is also forwarding to 1.1.1.1 (R1)

```
R4# show ip route 10.1.1.1

Routing entry for 10.0.0.0/8

Known via "bgp 65001", distance 200, metric 0

Routing Descriptor Blocks:

* 1.1.1.1, from 5.5.5.5, 01:47:02 ago

Route metric is 0, traffic share count is 1

AS Hops 0
```

Routing Loop - Diagnosis

R2 is forwarding to 3.3.3.3? (R3)

```
R2# show ip route 10.1.1.1

Routing entry for 10.0.0.0/8

Known via "bgp 65000", distance 200, metric 0

Routing Descriptor Blocks:

* 3.3.3.3, from 3.3.3.3, 01:47:00 ago

Route metric is 0, traffic share count is 1

AS Hops 0, BGP network version 3
```

 Very odd that the NEXT_HOP is in the middle of the network

Routing Loop - Diagnosis

Verify BGP paths on R2

```
R2#show ip bgp 10.0.0.0

BGP routing table entry for 10.0.0.0/8, version 3

Paths: (4 available, best #1)

Advertised to non peer-group peers:

1.1.1.1 5.5.5.5 4.4.4.4

(65001 65002)

3.3.3.3 (metric 11) from 3.3.3.3 (3.3.3.3)

Origin IGP, metric 0, localpref 100, valid, confedinternal, best

(65002)

1.1.1.1 (metric 5010) from 1.1.1.1 (1.1.1.1)

Origin IGP, metric 0, localpref 100, valid, confedexternal
```

- R3 path is better than R1 path because of IGP cost to NEXT_HOP
- R3 is advertising the path to us with a NEXT_HOP of 3.3.3.3 ????

Routing Loop - Diagnosis

• What is R3 advertising?

```
R3# show ip bgp 10.0.0.0

BGP routing table entry for 10.0.0.0/8, version 3

Paths: (2 available, best #1, table Default-IP-Routing-Table)

Advertised to non peer-group peers:

5.5.5.5 2.2.2.2

(65001 65002)

1.1.1.1 (metric 5031) from 4.4.4.4 (4.4.4.4)

Origin IGP, metric 0, localpref 100, valid, confedexternal, best, multipath

(65001 65002)

1.1.1.1 (metric 5031) from 5.5.5.5 (5.5.5.5)

Origin IGP, metric 0, localpref 100, valid, confedexternal, multipath
```

Hmmm, R3 is using multipath to load-balance

R3#show run | include maximum

Routing Loop - Solution

 "maximum-paths" tells the router to reset the NEXT_HOP to himself

R3 sets NEXT_HOP to 3.3.3.3

- Forces traffic to come to him so he can loadbalance
- Is typically used for multiple eBGP sessions to an AS

Be careful when using in Confederations!!

 Need to make R2 prefer the path from R1 to prevent the routing loop

Make IGP metric to 1.1.1.1 better than IGP metric to 4.4.4.4

- High CPU in "Router BGP" is normally a sign of a convergence problem
- Find a prefix that changes every minute show ip route | include, 00:00
- Troubleshoot/debug that one prefix

BGP routing loop?
 First, check for IGP routing loops to BGP NEXT_HOPs

 BGP loops are normally caused by Not following physical topology in RR environment Multipath within confederations Lack of a full iBGP mesh

Get the following from each router in the loop path

```
show ip route x.x.x.x show ip bgp x.x.x.x show ip route NEXT_HOP____
```

"show ip bgp neighbor x.x.x.x advertised-routes"

Lets you see a list of NLRI that you sent a peer

Note: The attribute values shown are taken from the BGP table. Attribute modifications by outbound routemaps will not be shown.

- "show ip bgp neighbor x.x.x.x routes"
 - Displays routes x.x.x.x sent to us that made it through our inbound filters
- "show ip bgp neighbor x.x.x.x received-routes"

 Can only use if "soft-reconfig inbound" is configured

 Displaye all routes received from a poor, even those

Displays all routes received from a peer, even those that were denied

- "clear ip bgp x.x.x.x in"
 Ask x.x.x.x to resend his UPDATEs to us
- "clear ip bgp x.x.x.x out"
 Tells BGP to resend UPDATEs to x.x.x.x
- "debug ip bgp update"
 Always use an ACL to limit output
 Great for troubleshooting "Automatic Denies"
- "debug ip bgp x.x.x.x update"
 Allows you to debug updates to/from a specific peer
 Handy if multiple peers are sending you the same prefix

Summary/Tips

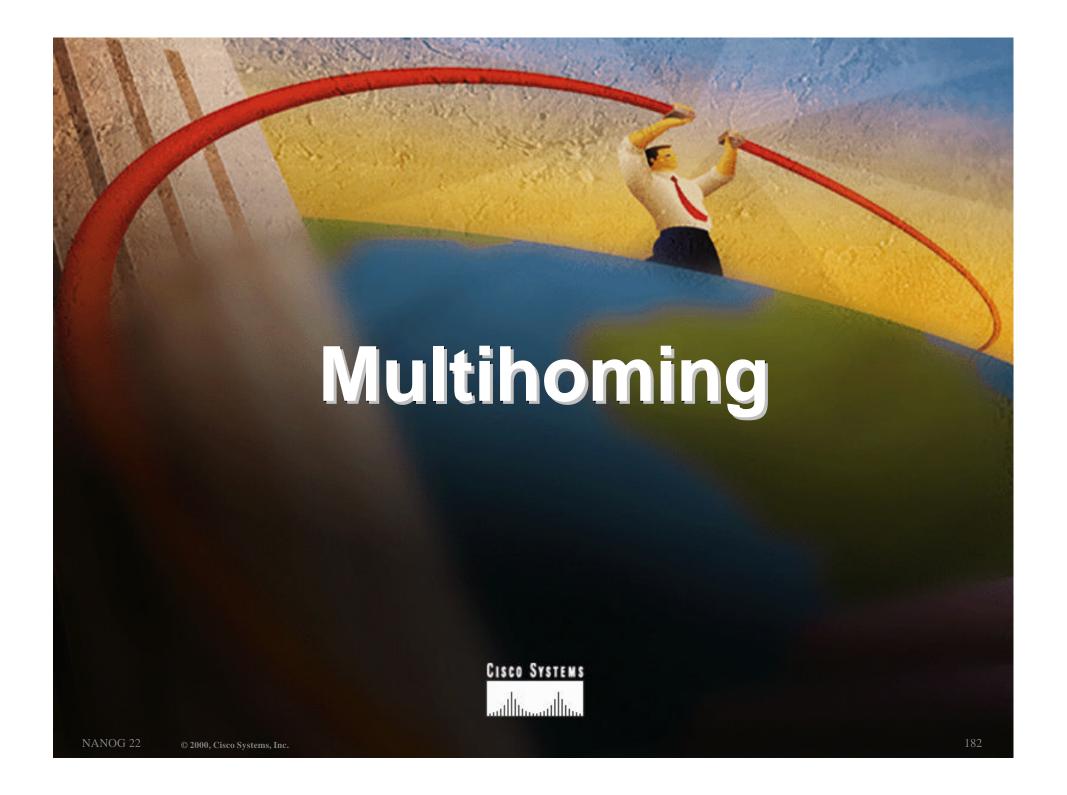
- Isolate the problem!!
- Use ACLs when enabling debug commands
- Enable bgp log-neighbor-changes
- IP reachability must exist for sessions to be established

learned from IGP

make sure the source and destination addresses match the configuration

BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities



Multihoming Definition

- More than one link external to the local network
 - two or more links to the same ISP two or more links to different ISPs
- Usually two external facing routers one router gives link and provider redundancy only

AS Numbers

- An Autonomous System Number is required by BGP
- Obtained from upstream ISP or Regional Registry
- Necessary when you have links to more than one ISP or exchange point

Configuring Policy

- Three BASIC Principles
 prefix-lists to filter prefixes
 filter-lists to filter ASNs
 route-maps to apply policy
- Avoids confusion!

Originating Prefixes

Basic Assumptions

MUST announce assigned address block to Internet

MAY also announce subprefixes – reachability is not guaranteed

RIR minimum allocation is /20

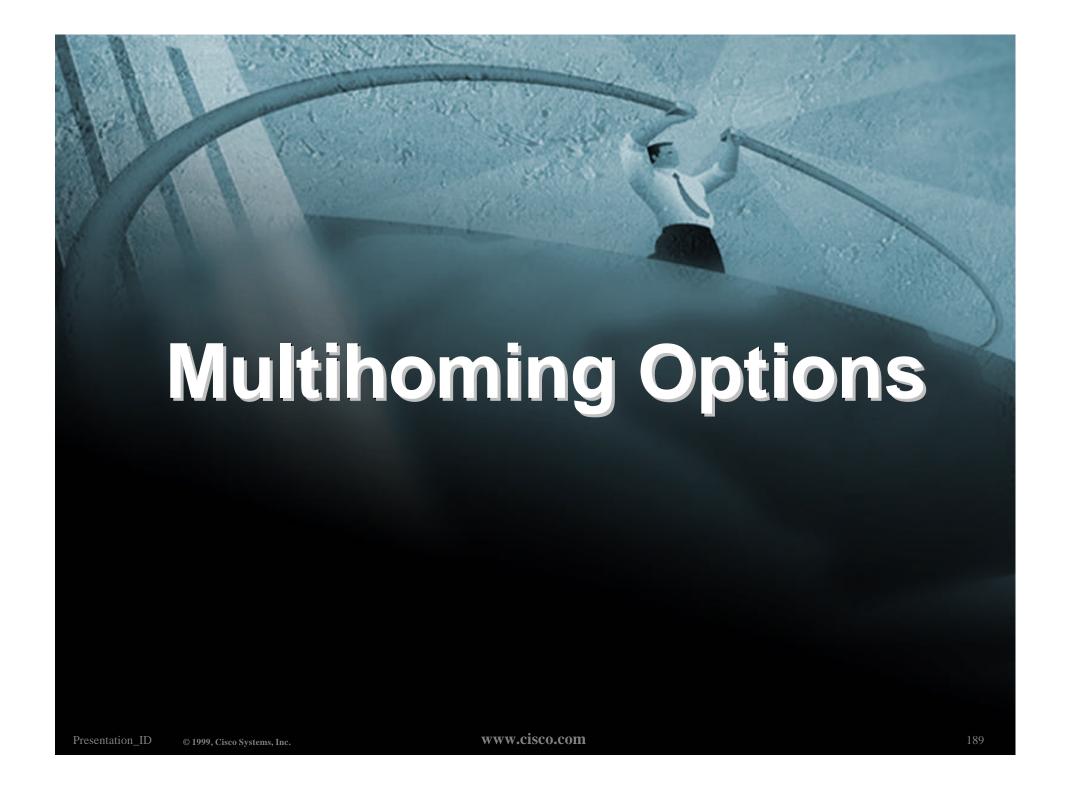
several ISPs filter RIR blocks on this boundary called "Net Police" by some

Part of the "Net Police" prefix list

```
!! APNIC
ip prefix-list FILTER permit 61.0.0.0/8 ge 9 le 20
ip prefix-list FILTER permit 202.0.0.0/7 ge 9 le 20
ip prefix-list FILTER permit 210.0.0.0/7 ge 9 le 20
ip prefix-list FILTER permit 218.0.0.0/8 ge 9 le 20
!! ARIN
ip prefix-list FILTER permit 63.0.0.0/8 ge 9 le 20
ip prefix-list FILTER permit 64.0.0.0/7 ge 9 le 20
ip prefix-list FILTER permit 66.0.0.0/8 ge 9 le 20
ip prefix-list FILTER permit 199.0.0.0/8 ge 9 le 20
ip prefix-list FILTER permit 200.0.0.0/8 ge 9 le 20
ip prefix-list FILTER permit 204.0.0.0/6 ge 9 le 20
ip prefix-list FILTER permit 208.0.0.0/7 ge 9 le 20
ip prefix-list FILTER permit 216.0.0.0/8 ge 9 le 20
!! RIPE NCC
ip prefix-list FILTER permit 62.0.0.0/8 ge 9 le 20
ip prefix-list FILTER permit 80.0.0.0/7 ge 9 le 20
ip prefix-list FILTER permit 193.0.0.0/8 ge 9 le 20
ip prefix-list FILTER permit 194.0.0.0/7 ge 9 le 20
ip prefix-list FILTER permit 212.0.0.0/7 ge 9 le 20
```

"Net Police" prefix list issues

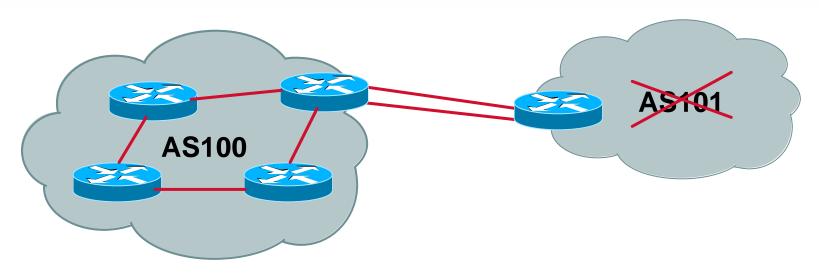
- meant to "punish" ISPs who won't and don't aggregate
- impacts legitimate multihoming
- impacts regions where domestic backbone is unavailable or costs \$\$\$ compared with international bandwidth
- hard to maintain requires updating when RIRs start allocating from new address blocks
- don't do it unless consequences understood and you are prepared to keep it current



Multihoming Scenarios

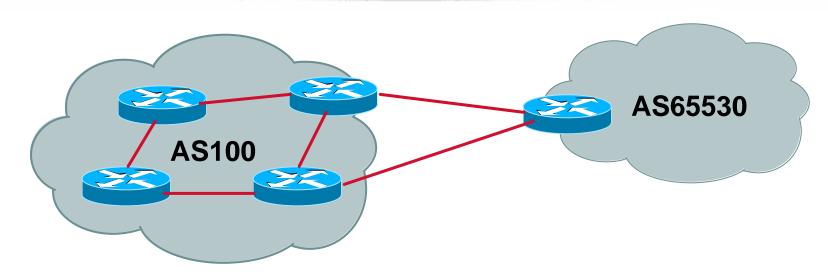
- Stub network
- Multi-homed stub network
- Multi-homed network
- Configuration Options

Stub Network



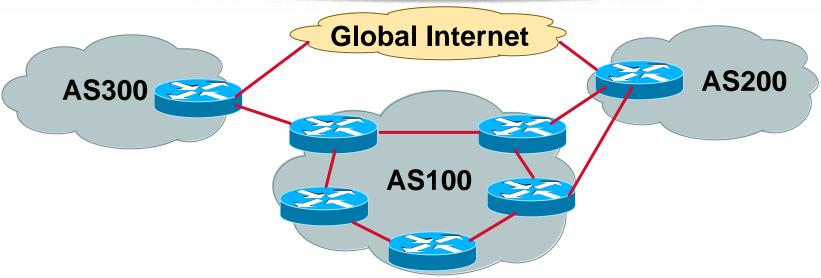
- No need for BGP
- Point static default to upstream ISP
- Upstream ISP advertises stub network
- Policy confined within upstream ISP's policy

Multi-homed Stub Network



- Use BGP (not IGP or static) to loadshare
- Use private AS (ASN > 64511)
- Upstream ISP advertises stub network
- Policy confined within upstream ISP's policy

Multi-Homed Network

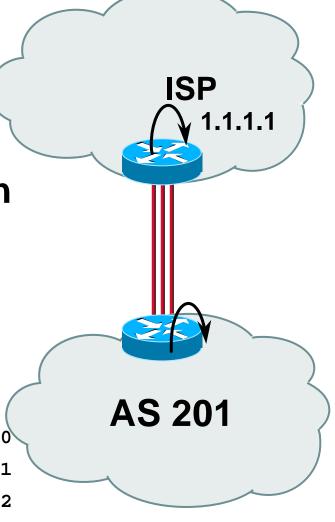


Many situations possible
 multiple sessions to same ISP
 secondary for backup only
 load-share between primary and secondary
 selectively use different ISPs

Multiple Sessions to an ISP - Example One

- eBGP multihop
- eBGP to loopback addresses
- eBGP prefixes learned with loopback address as next hop

```
router bgp 201
neighbor 1.1.1.1 remote-as 200
neighbor 1.1.1.1 ebgp-multihop 5
ip route 1.1.1.1 255.255.255.255 serial 1/0
ip route 1.1.1.1 255.255.255.255 serial 1/1
ip route 1.1.1.1 255.255.255.255 serial 1/2
```



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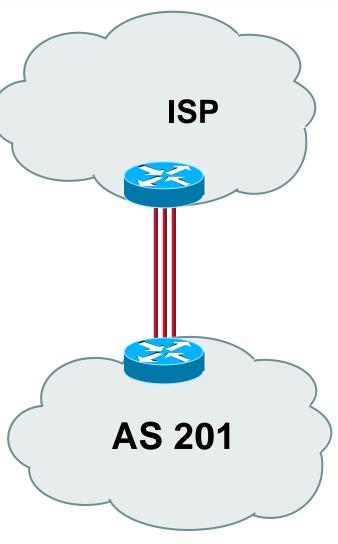
Multiple Sessions to an ISP - Example Two

BGP multi-path

NANOG 22

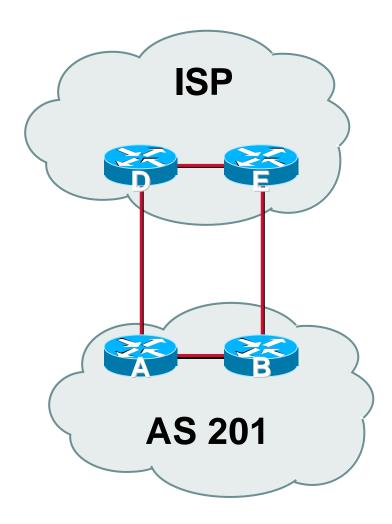
- Three BGP sessions required
- limit of 6 parallel paths

```
router bgp 201
neighbor 1.1.2.1 remote-as 200
neighbor 1.1.2.5 remote-as 200
neighbor 1.1.2.9 remote-as 200
maximum-paths 3
```



Multiple Sessions to an ISP

- Simplest scheme is to use defaults
- Learn/advertise prefixes for better control
- Planning and some work required to achieve loadsharing
- No magic solution

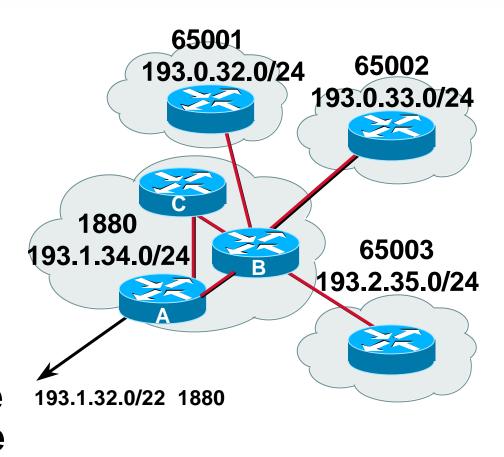


Private-AS – Application

Applications

ISP with singlehomed customers (RFC2270)

corporate network with several regions and connections to the Internet only in the core



Private-AS Removal

- neighbor x.x.x.x remove-private-AS
- Rules:

available for eBGP neighbors only

if the update has AS_PATH made up of private-AS numbers, the private-AS will be dropped

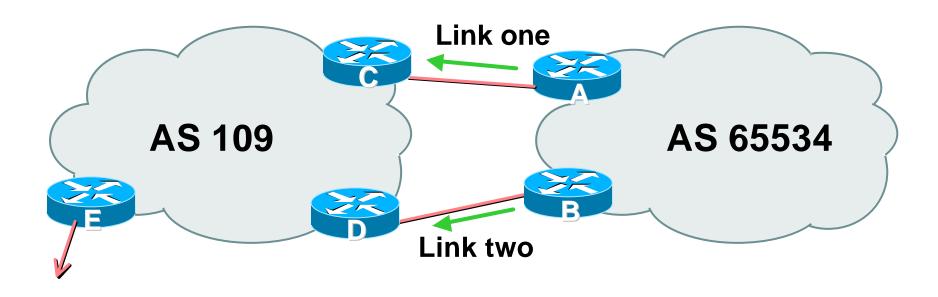
if the AS_PATH includes private and public AS numbers, private AS number will not be removed...it is a configuration error!

if AS_PATH contains the AS number of the eBGP neighbor, the private-AS numbers will not be removed

if used with confederations, it will work as long as the private AS numbers are after the confederation portion of the AS_PATH

Two links to the same **ISP** With Redundancy and Loadsharing

Two links to the same ISP (with redundancy)



 AS109 removes private AS and any customer subprefixes from Internet announcement

Loadsharing to the same ISP

- Announce /19 aggregate on each link
- Split /19 and announce as two /20s, one on each link basic inbound loadsharing assumes equal circuit capacity and even spread of traffic across address block
- Vary the split until "perfect" loadsharing achieved
- Accept the default from upstream
 basic outbound loadsharing by nearest exit
 okay in first approx as most ISP and end-site traffic is inbound

Two links to the same ISP

Router A Configuration

```
router bgp 65534
network 221.10.0.0 mask 255.255.224.0
network 221.10.0.0 mask 255.255.240.0
neighbor 222.222.10.2 remote-as 109
neighbor 222.222.10.2 prefix-list routerC out
neighbor 222.222.10.2 prefix-list default in
ip prefix-list default permit 0.0.0.0/0
ip prefix-list routerC permit 221.10.0.0/20
ip prefix-list routerC permit 221.10.0.0/19
ip route 221.10.0.0 255.255.240.0 null0
ip route 221.10.0.0 255.255.224.0 null0
```

Router B configuration is similar but with the other /20

Two links to the same ISP

Router C Configuration

```
router bgp 109
neighbor 222.222.10.1 remote-as 65534
neighbor 222.222.10.1 default-originate
neighbor 222.222.10.1 prefix-list Customer in
neighbor 222.222.10.1 prefix-list default out
!
ip prefix-list Customer permit 221.10.0.0/19 le 20
ip prefix-list default permit 0.0.0.0/0
```

- Router C only allows in /19 and /20 prefixes from customer block
- Router D configuration is identical

Loadsharing to the same ISP

- Loadsharing configuration is only on customer router
- Upstream ISP has to

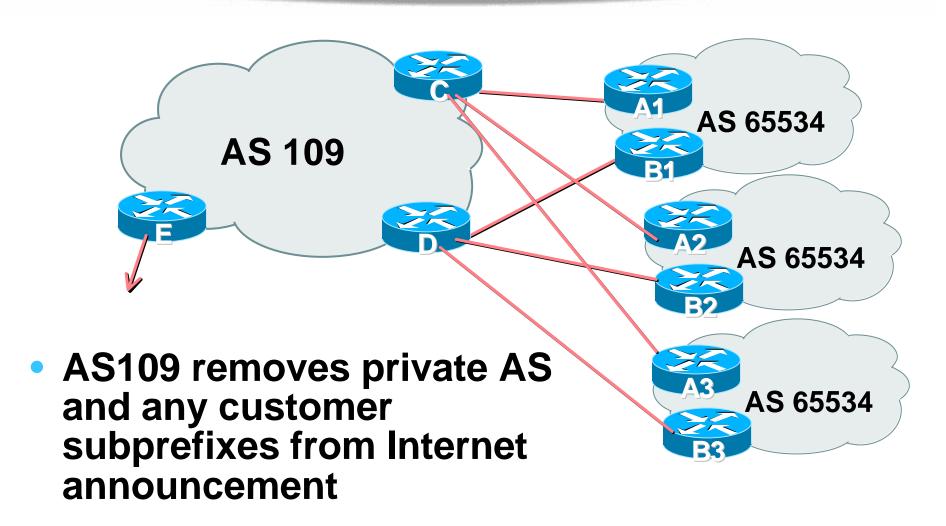
remove customer subprefixes from external announcements

remove private AS from external announcements

Could also use BGP communities



Multiple Dualhomed Customers (RFC2270)



- Customer announcements as per previous example
- Use the same private AS for each customer documented in RFC2270 address space is not overlapping each customer hears default only
- Router An and Bn configuration same as Router A and B previously

Two links to the same ISP

Router A1 Configuration

```
router bgp 65534
network 221.10.0.0 mask 255.255.224.0
network 221.10.0.0 mask 255.255.240.0
neighbor 222.222.10.2 remote-as 109
neighbor 222.222.10.2 prefix-list routerC out
neighbor 222.222.10.2 prefix-list default in
ip prefix-list default permit 0.0.0.0/0
ip prefix-list routerC permit 221.10.0.0/20
ip prefix-list routerC permit 221.10.0.0/19
ip route 221.10.0.0 255.255.240.0 null0
ip route 221.10.0.0 255.255.224.0 null0
```

Router B1 configuration is similar but for the other /20

Router C Configuration

```
router bgp 109
neighbor bgp-customers peer-group
neighbor bgp-customers remote-as 65534
neighbor bgp-customers default-originate
neighbor bgp-customers prefix-list default out
neighbor 222.222.10.1 peer-group bgp-customers
neighbor 222.222.10.1 description Customer One
neighbor 222.222.10.1 prefix-list Customer1 in
neighbor 222.222.10.9 peer-group bgp-customers
neighbor 222.222.10.9 description Customer Two
neighbor 222.222.10.9 prefix-list Customer2 in
```

```
neighbor 222.222.10.17 peer-group bgp-customers
neighbor 222.222.10.17 description Customer Three
neighbor 222.222.10.17 prefix-list Customer3 in
!
ip prefix-list Customer1 permit 221.10.0.0/19 le 20
ip prefix-list Customer2 permit 221.16.64.0/19 le 20
ip prefix-list Customer3 permit 221.14.192.0/19 le 20
ip prefix-list default permit 0.0.0.0/0
```

- Router C only allows in /19 and /20 prefixes from customer block
- Router D configuration is almost identical

Router E Configuration

assumes customer address space is not part of upstream's address block

```
router bgp 109
neighbor 222.222.10.17 remote-as 110
neighbor 222.222.10.17 remove-private-AS
neighbor 222.222.10.17 prefix-list Customers out
!
ip prefix-list Customers permit 221.10.0.0/19
ip prefix-list Customers permit 221.16.64.0/19
ip prefix-list Customers permit 221.14.192.0/19
```

Private AS still visible inside AS109

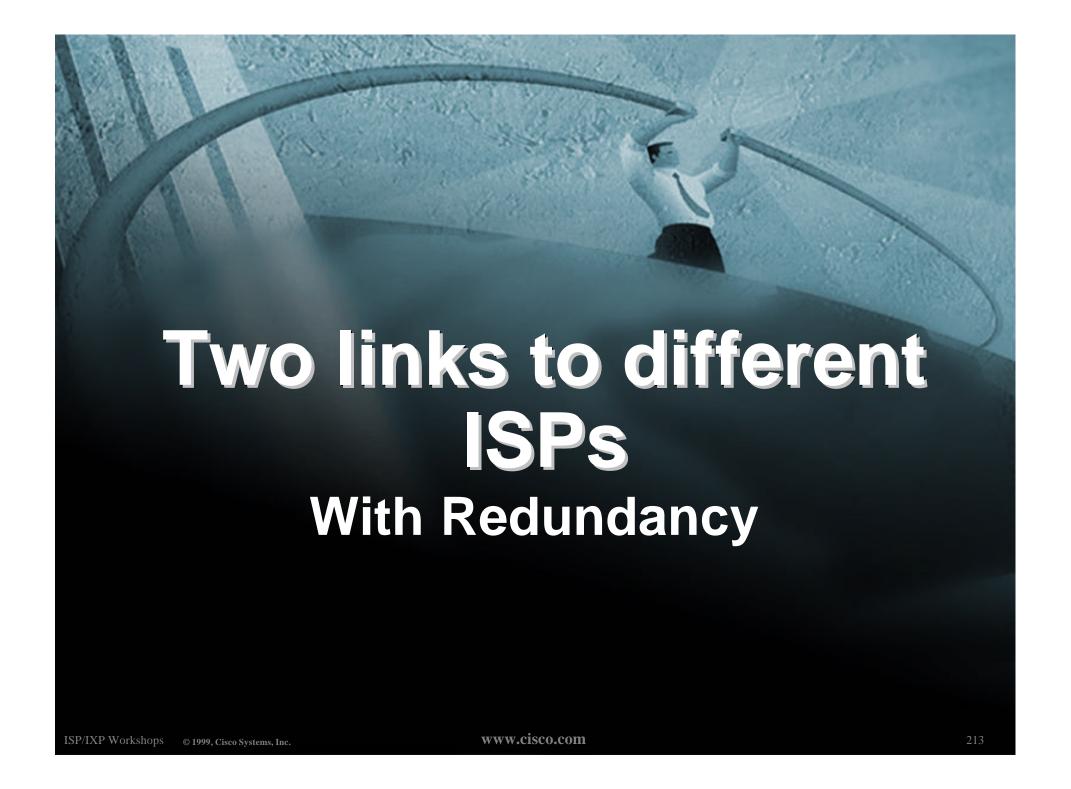
 If customers' prefixes come from ISP's address block

do NOT announce them to the Internet

announce ISP aggregate only

Router E configuration:

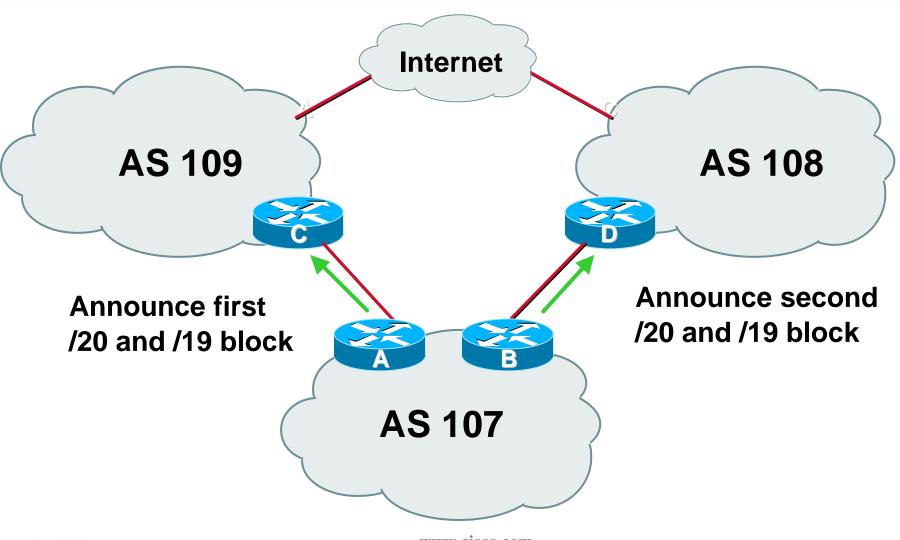
```
router bgp 109
neighbor 222.222.10.17 remote-as 110
neighbor 222.222.10.17 prefix-list my-aggregate out
!
ip prefix-list my-aggregate permit 221.8.0.0/13
```



Two links to different ISPs (with redundancy)

- Announce /19 aggregate on each link
- Split /19 and announce as two /20s, one on each link
 - basic inbound loadsharing
- When one link fails, the announcement of the /19 aggregate via the other ISP ensures continued connectivity

Two links to different ISPs (with redundancy)



Two links to different ISPs (with redundancy)

Router A Configuration

```
router bgp 107
network 221.10.0.0 mask 255.255.224.0
network 221.10.0.0 mask 255.255.240.0
neighbor 222.222.10.1 remote-as 109
neighbor 222.222.10.1 prefix-list firstblock out
neighbor 222.222.10.1 prefix-list default in
ip prefix-list default permit 0.0.0.0/0
ip prefix-list firstblock permit 221.10.0.0/20
ip prefix-list firstblock permit 221.10.0.0/19
```

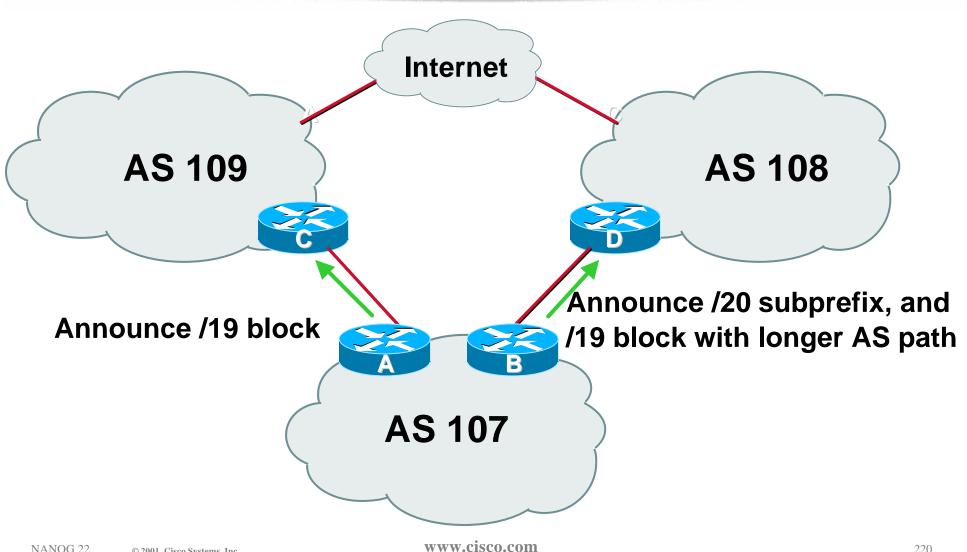
Two links to different ISPs (with redundancy)

Router B Configuration

```
router bgp 107
 network 221.10.0.0 mask 255.255.224.0
network 221.10.16.0 mask 255.255.240.0
neighbor 220.1.5.1 remote-as 108
neighbor 220.1.5.1 prefix-list secondblock out
neighbor 220.1.5.1 prefix-list default in
ip prefix-list default permit 0.0.0.0/0
ip prefix-list secondblock permit 221.10.16.0/20
ip prefix-list secondblock permit 221.10.0.0/19
```

Two links to different ISPs **More Controlled Loadsharing**

- Announce /19 aggregate on each link
 - On first link, announce /19 as normal
 - On second link, announce /19 with longer AS PATH, and announce one /20 subprefix
 - controls loadsharing between upstreams and the Internet
- Vary the subprefix size and AS PATH length until "perfect" loadsharing achieved
- Still require redundancy!



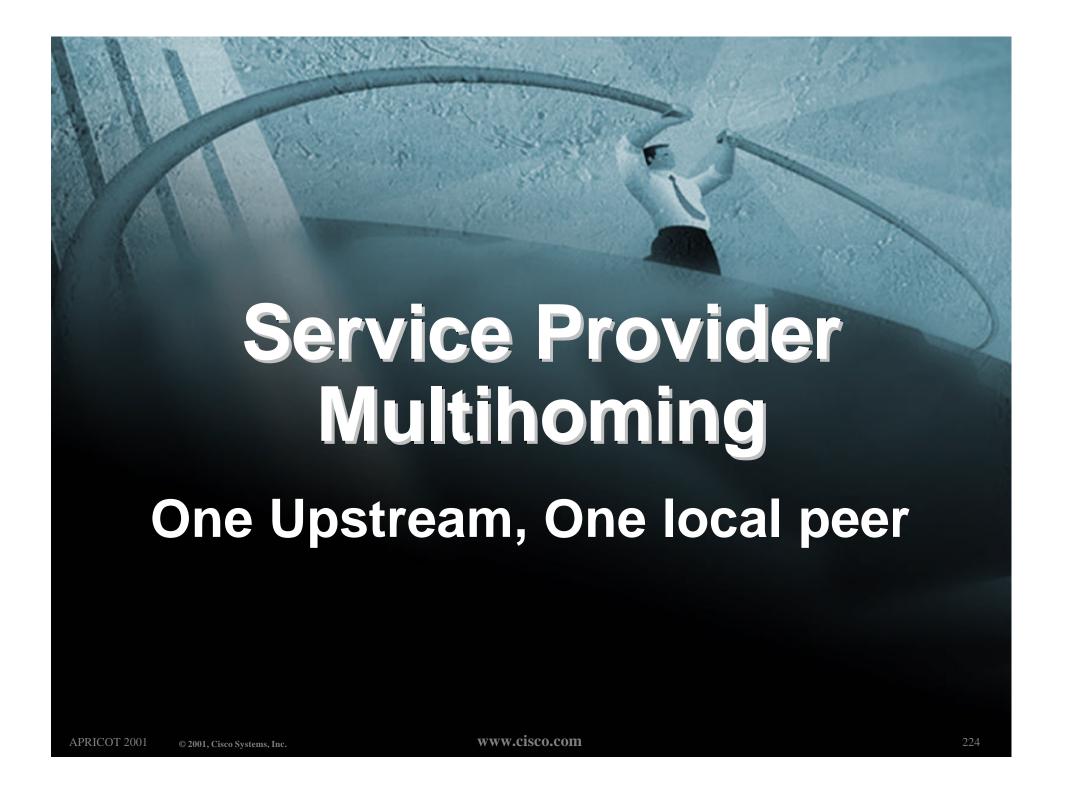
Router A Configuration

```
router bgp 107
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 109
neighbor 222.222.10.1 prefix-list default in
neighbor 222.222.10.1 prefix-list aggregate out
!
ip prefix-list aggregate permit 221.10.0.0/19
```

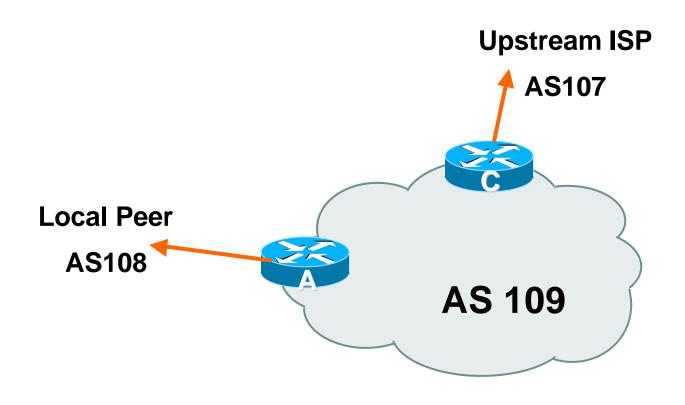
Router B Configuration

```
network 221.10.0.0 mask 255.255.224.0
network 221.10.16.0 mask 255.255.240.0
neighbor 220.1.5.1 remote-as 108
neighbor 220.1.5.1 prefix-list default in
neighbor 220.1.5.1 prefix-list subblocks out
neighbor 220.1.5.1 route-map routerD out
!
..next slide..
```

```
route-map routerD permit 10
match ip address prefix-list aggregate
set as-path prepend 107 107
route-map routerD permit 20
!
ip prefix-list subblocks permit 221.10.0.0/19 le 20
ip prefix-list aggregate permit 221.10.0.0/19
```



- Announce /19 aggregate on each link
- Accept default route only from upstream
 Either 0.0.0.0/0 or a network which can be used as default
- Accept all routes from local peer
- Border routers talk iBGP with each other



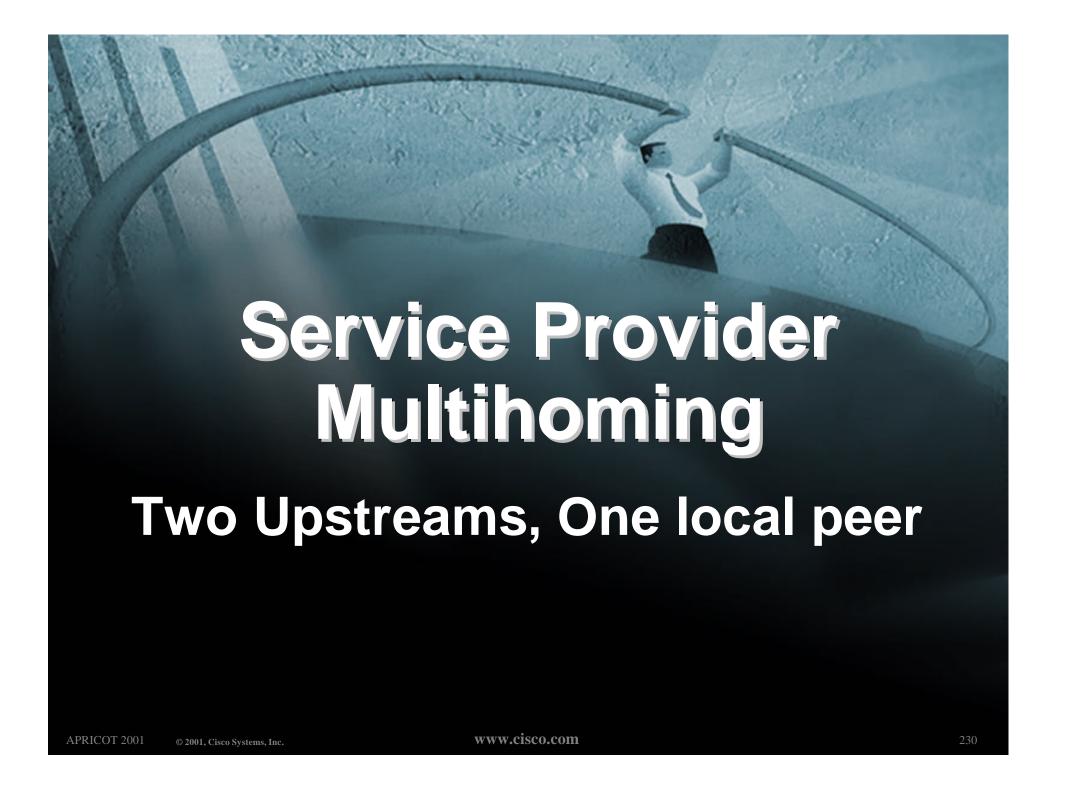
Router A Configuration

```
router bgp 109
 network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.2 remote-as 108
neighbor 222.222.10.2 prefix-list my-block out
neighbor 222.222.10.2 prefix-list AS108-peer in
ip prefix-list AS108-peer permit 222.5.16.0/19
ip prefix-list AS108-peer permit 221.240.0.0/20
ip prefix-list my-block permit 221.10.0.0/19
ip route 221.10.0.0 255.255.224.0 null0
```

Router C Configuration

```
router bgp 109
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 107
neighbor 222.222.10.1 prefix-list default in
neighbor 222.222.10.1 prefix-list my-block out
ip prefix-list my-block permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
ip route 221.10.0.0 255.255.224.0 null0
```

- Two configurations possible for Router A
 - Filtering on ASes assumes peer knows what they are doing (never do this)
 - Prefix-list higher maintenance, but safer
- Local traffic goes to and from local peer, everything else goes to upstream
- Routers A and C have minimum memory and CPU requirements



Two Upstreams, One Local Peer

• Two configuration options:

Accept full routing from both upstreams Expensive!

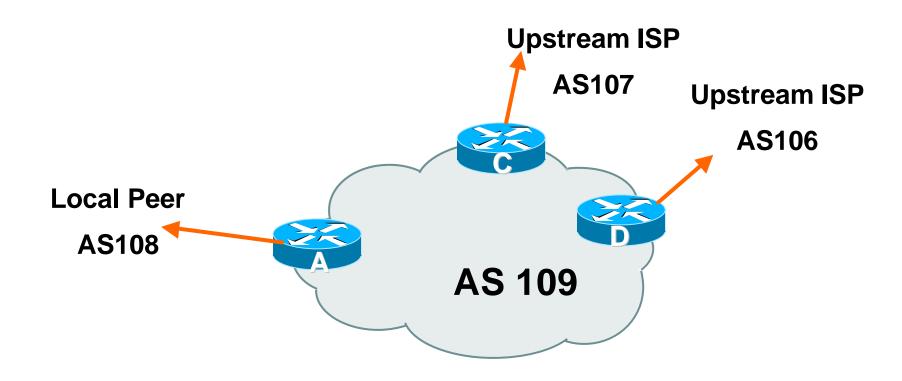
But this is the popular choice today?!!

Accept default from one upstream and some routes from the other upstream

Best compromise, not expensive!

Better convergence rate and stability

Two Upstreams, One Local Peer



Router A configuration is as previously

Router C Configuration

```
router bgp 109
 network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 107
neighbor 222.222.10.1 prefix-list rfc1918-deny in
neighbor 222.222.10.1 prefix-list my-block out
neighbor 222.222.10.1 route-map AS107-loadshare in
ip prefix-list my-block permit 221.10.0.0/19
! See earlier in tutorial for RFC1918 list
ip route 221.10.0.0 255.255.224.0 null0
..next slide
```

```
ip as-path access-list 10 permit ^(107_)+$
ip as-path access-list 10 permit ^(107_)+_[0-9]+$
!
route-map AS107-loadshare permit 10
  match ip as-path 10
  set local-preference 120
route-map AS107-loadshare permit 20
  set local-preference 80
!
```

Router C configuration:

Accept full routes from AS107

Tag prefixes originated by AS107 and AS107's neighbouring ASes with local preference 120

Remaining prefixes tagged with local preference of 80

Traffic to those ASes will go over AS107 link

Traffic to other all other ASes will go over the link to AS106

Router D configuration same as Router C without the route-map

Hears full routing table!

Full routes from upstreams

Expensive – needs lots of memory today

Expensive – contributes to network instability

Need to play preference games

Previous example is only an example – real life will need improved fine-tuning!

Previous example doesn't consider inbound traffic – see earlier slides for examples

Router C Configuration

```
router bgp 109
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 107
neighbor 222.222.10.1 prefix-list rfc1918-nodef-deny in
neighbor 222.222.10.1 prefix-list my-block out
neighbor 222.222.10.1 filter-list 10 in
neighbor 222.222.10.1 route-map tag-default-low in
ip prefix-list my-block permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
! See earlier in tutorial for RFC1918 list
ip route 221.10.0.0 255.255.224.0 null0
```

```
ip as-path access-list 10 permit ^(107_)+$
ip as-path access-list 10 permit ^(107_)+_[0-9]+$
!
route-map tag-default-low permit 10
match ip address prefix-list default
set local-preference 80
route-map tag-default-low permit 20
!
```

Router D Configuration

```
router bgp 109
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.5 remote-as 106
neighbor 222.222.10.5 prefix-list default in
neighbor 222.222.10.5 prefix-list my-block out
ip prefix-list my-block permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
ip route 221.10.0.0 255.255.224.0 null0
```

Router C configuration:

Accept full routes from AS107

(or get them to send less)

Filter ASNs so only AS107 and AS107's neighbouring ASes are accepted

Allow default, and set it to local preference 80

Traffic to those ASes will go over AS107 link

Traffic to other all other ASes will go over the link to AS106

If AS106 link fails, backup via AS107 – and viceversa

Partial routes from upstreams

Not expensive – only carry the routes necessary for loadsharing

Not expensive – network more stable!

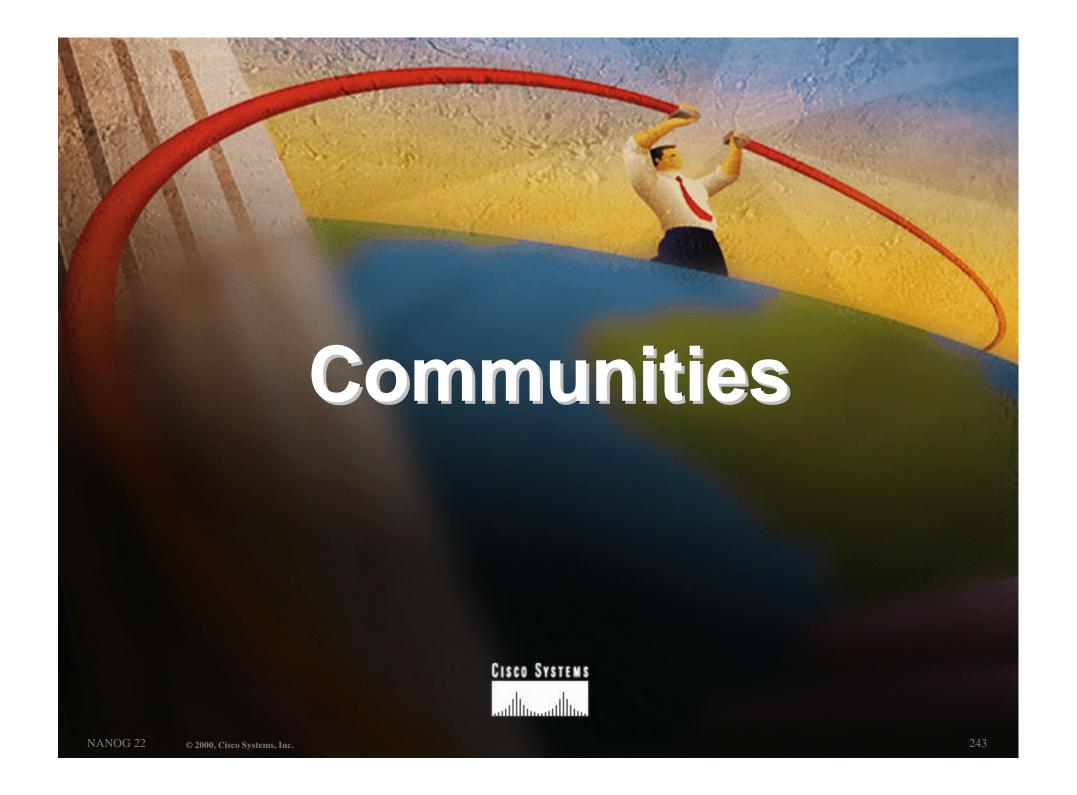
Need to filter on AS paths

Previous example is only an example – real life will need improved fine-tuning!

Previous example doesn't consider inbound traffic – see earlier slides for examples

BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities



Community usage

- RFC1998
- Examples of SP applications

- Informational RFC
- Describes how to implement loadsharing and backup on multiple inter-AS links
 - BGP communities used to determine local preference in upstream's network
- Gives control to the customer
- Simplifies upstream's configuration simplifies network operation!

Community values defined to have particular meanings:

ASx:100 set local pref 100 preferred route

ASx:90 set local pref 90 backup route if dualhomed on ASx

ASx:80 set local pref 80 main link is to another ISP with

same AS path length

ASx:70 set local pref 70 main link is to another ISP

Sample Customer Router Configuration

```
router bgp 107
neighbor x.x.x.x remote-as 109
neighbor x.x.x.x description Backup ISP
neighbor x.x.x.x route-map config-community out
neighbor x.x.x.x send-community
ip as-path access-list 20 permit ^$
ip as-path access-list 20 deny .*
route-map config-community permit 10
match as-path 20
 set community 109:90
```

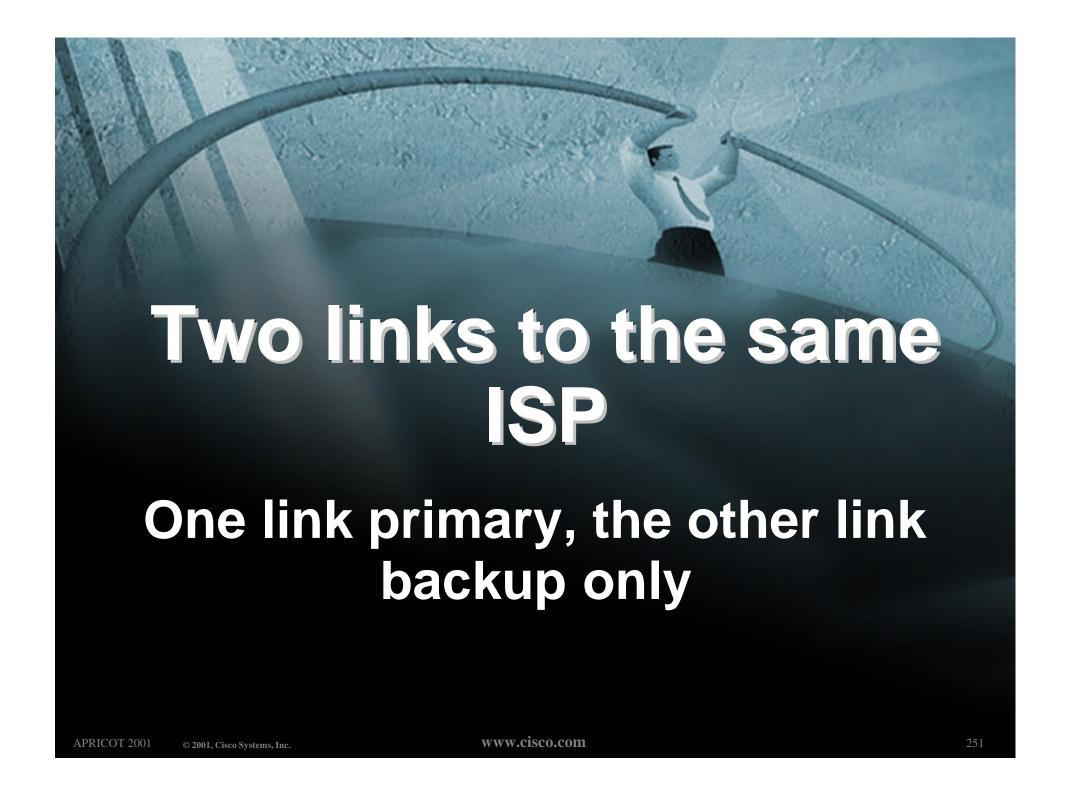
Sample ISP Router Configuration

```
! Homed to another ISP
ip community-list 70 permit 109:70
! Homed to another ISP with equal ASPATH length
ip community-list 80 permit 109:80
! Customer backup routes
ip community-list 90 permit 109:90
route-map set-customer-local-pref permit 10
match community 70
 set local-preference 70
```

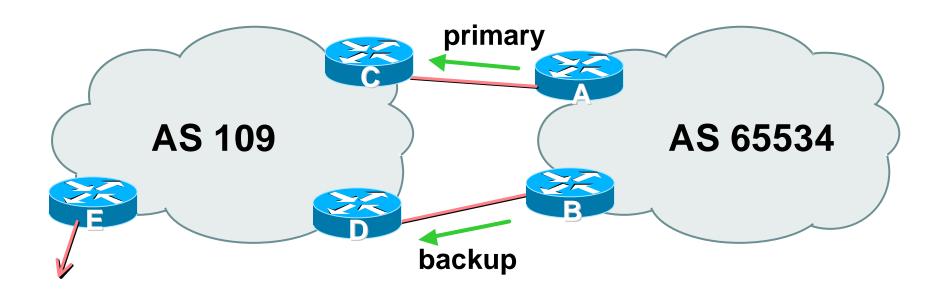
Sample ISP Router Configuration

```
route-map set-customer-local-pref permit 20
match community 80
 set local-preference 80
route-map set-customer-local-pref permit 30
match community 90
 set local-preference 90
route-map set-customer-local-pref permit 40
 set local-preference 100
```

Supporting RFC1998
 many ISPs do, more should
 check AS object in the Internet
 Routing Registry
 if you do, insert comment in AS object
 in the IRR



Two links to the same ISP



AS109 proxy aggregates for AS 65534

- Announce /19 aggregate on each link primary link makes standard announcement
 - backup link sends community
- When one link fails, the announcement of the /19 aggregate via the other link ensures continued connectivity

Router A Configuration

```
router bgp 65534
 network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.2 remote-as 109
neighbor 222.222.10.2 description RouterC
neighbor 222.222.10.2 prefix-list aggregate out
 neighbor 222.222.10.2 prefix-list default in
ip prefix-list aggregate permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
```

Router B Configuration

```
router bgp 65534
 network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.6 remote-as 109
 neighbor 222.222.10.6 description RouterD
 neighbor 222.222.10.6 send-community
neighbor 222.222.10.6 prefix-list aggregate out
neighbor 222.222.10.6 route-map routerD-out out
 neighbor 222.222.10.6 prefix-list default in
neighbor 222.222.10.6 route-map routerD-in in
.. next slide
```

```
ip prefix-list aggregate permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
route-map routerD-out permit 10
match ip address prefix-list aggregate
 set community 109:90
route-map routerD-out permit 20
route-map routerD-in permit 10
 set local-preference 90
```

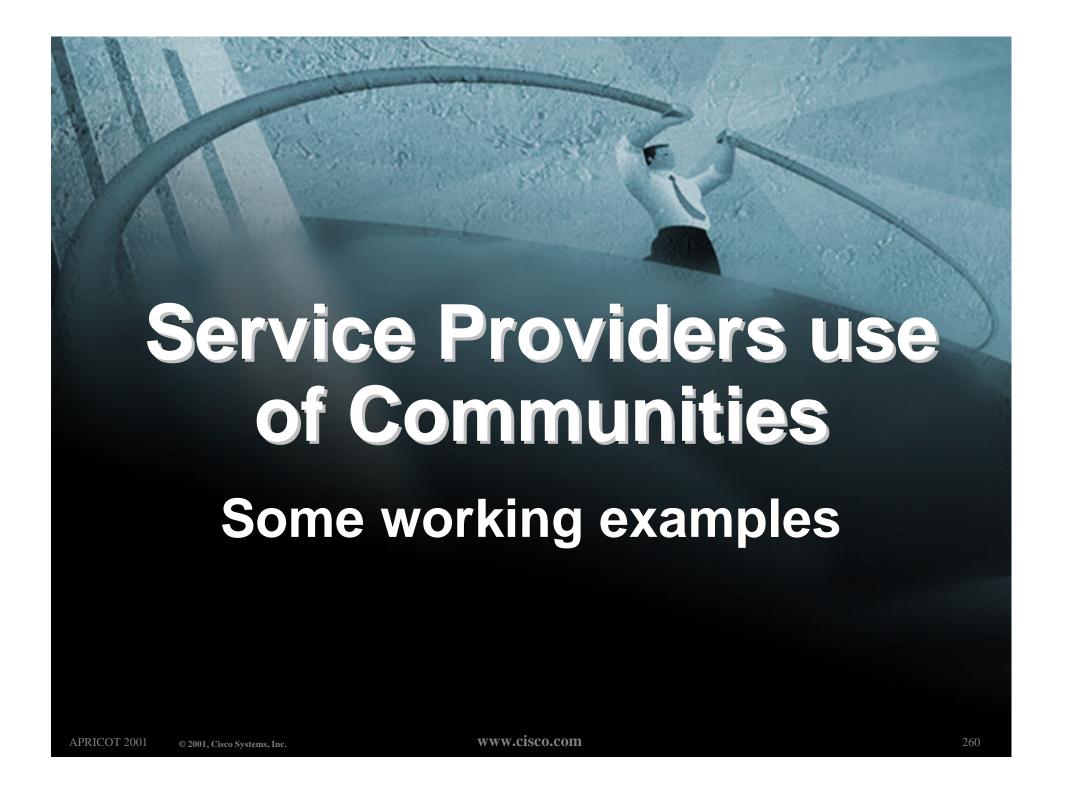
Router C Configuration (main link)

```
router bgp 109
neighbor 222.222.10.1 remote-as 65534
neighbor 222.222.10.1 default-originate
neighbor 222.222.10.1 prefix-list Customer in
neighbor 222.222.10.1 prefix-list default out
!
ip prefix-list Customer permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
```

Router D Configuration (backup link)

```
router bgp 109
neighbor 222.222.10.5 remote-as 65534
neighbor 222.222.10.5 default-originate
neighbor 222.222.10.5 prefix-list Customer in
neighbor 222.222.10.5 route-map bgp-cust-in in
neighbor 222.222.10.5 prefix-list default out
ip prefix-list Customer permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
..next slide
```

```
ip prefix-list Customer permit 221.10.0.0/19
  ip prefix-list default permit 0.0.0.0/0
  ip community-list 90 permit 109:90
<snip>
  route-map bgp-cust-in permit 30
  match community 90
   set local-preference 90
  route-map bgp-cust-in permit 40
   set local-preference 100
```



Background

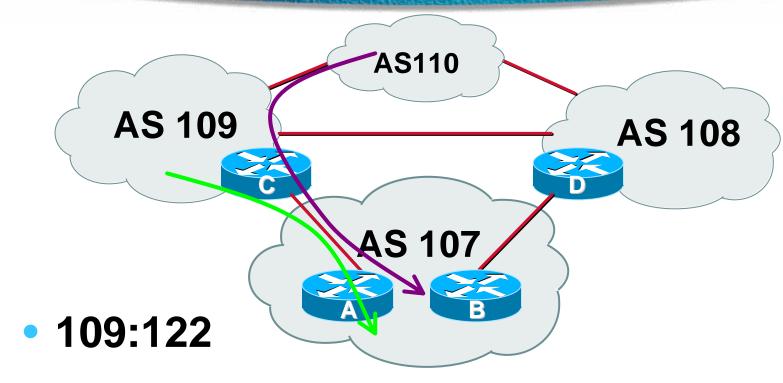
- RFC1998 is okay for "simple" multihomed customers
 - assumes that upstreams are interconnected
- ISPs create many other communities to handle more complex situations

More community definitions

```
ASx:122 set local pref 120 and set local pref high on upstreams
         set local pref 120 and set local pref low on upstreams
         set local pref 120 (opposite to ASx:80)
ASx:120
         set local pref 80 and set local pref high on upstreams
ASx:82
ASx:81
         set local pref 80 and set local pref low on upstreams
         announce to customers with no-export
ASx:21
ASx:20
         announce only to backbone and customers
ASx:3
         set 3x as-path prepend on peer announcement
ASx:2
         set 2x as-path prepend on peer announcement
ASx:1
         set 1x as-path prepend on peer announcement
(and variations on this theme depending on local conditions, e.g.
IXPs, domestic vs. international transit, etc.)
```

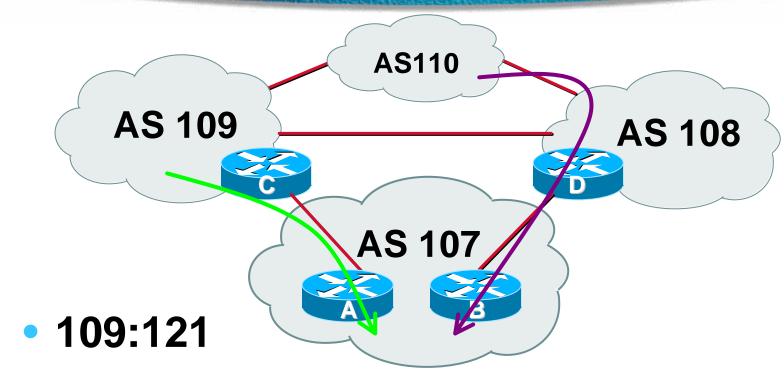
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Examples



traffic in AS109 comes directly to you traffic in AS110 sent to AS109 rather than best path

Examples



traffic in AS109 comes directly to you traffic in AS110 sent to AS108 rather than best path

Examples

109:3

prepend any announcements to peers of AS109 with 109_109_109

"AS109 is my backup transit AS"

109:20

Don't announce outside upstream's customer base

"AS109 provides local connections only"

109:21 is very similar

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BGP for Internet Service Providers **End of Tutorial** CISCO SYSTEMS NANOG 22 © 2000, Cisco Systems, Inc.